

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: AC Print, Octal Number, Sign, Single Column			LSR # OT 1.1 t
			Classification Closed
No. of Regs. in Subroutine	Temp. Regs. used by Subroutine	Average Time (operations)	Max. Time (operations)
35	1t, 2t	61.5	63.0
Preset Parameters none			
Program Parameters on entering Subroutine ac: Number to be printed			
Results on leaving Subroutine ac: 0			
Description <p>When control is transferred to this routine, the number to be printed should be in the accumulator. A five digit octal number is printed preceded by either a "0" for a positive or a "1" for a negative number. After the number is printed a carriage return is executed so that continual use of this routine will result in a single column of numbers.</p> <p>Notes:</p> <ol style="list-style-type: none"> 1. No point is printed. 2. The qp orders are set to suppress the punch. 			

MD, Nov.9,1951 JWC III, Nov.18,1951
--

Abstract: This subroutine prints a five digit octal number with a "0" or "1" and no point preceding it and then a carriage return. A "0" designates a positive number and a "1" a negative number.

Upon entering the subroutine:

AC: number to be printed

Temporary Registers:

1t: register used to store the remainder of the word.

2t: register used to store the digit counter.

00	ta 19r	Set return address	17	ca 4r	} Yes. Cause a carriage return
01	ts 1t	Store value	18	qp 128	
02	cp 20r	Is word negative?	19	<u>(sp 0)</u>	} Go back to main program
03	ca 26r	No. Print "0"	2r → 20	ad 24r	
23r → 04	qp 144	Printing	21	ts 1t	Store value
05	cs 25r	} Set digit counter	22	ca 27r	Print "1"
06	ts 2t		23	<u>sp 4r</u>	
16r → 07	ca 1t	Value in AC	24	0.77777	
08	sr*12		25	p4	Initial value of counter
09	ad 34r	} Add start of number table	26	p45	} Number Table
10	td 13r		27	p36	
11	sl 15	} Store	28	p39	
12	ts 1t	} remainder	29	p3	
13	(ca 0)	} Put flexo code for digit in AC	30	p21	
14	qp 128	Print digit	31	p33	
15	ao 2t	} Have all digits been printed? No.	32	p43	
16	cp 7r		33	p15	
			34	p26r	

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: AC Print and/or Punch, Octal Number (Magnitude), Sign, No Carriage Return			LSR # OT 1.2 t
			Classification Closed
No. of Regs. 33	Temp. Regs. 3	Average Time (operations) 7 print	Max. Time (operations) 7print
Preset Parameters v1 p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously			
Program Parameters ac: Number to be recorded			
Results ac: Number has been recorded and is no longer available, C(AC) = p1			
Description When control is transferred to this subroutine, the number to be printed should be in the accumulator. A five digit octal number is printed preceded by either a "+" for a positive number or a "-" for a negative number, i.e. 0.12345 → +12345 and 1.12345 → -65432.			
Notes: 1. No point is printed 2. No carriage return or space is printed 3. The <u>gp</u> instructions are normally set to print and suppress the punch.			

TITLE: AC PRINT AND/OR PUNCH, OCTAL NUMBER (MAGNITUDE), ISR# OT 1.2 t
SIGN, NO CARRIAGE RETURN

Abstract: This subroutine records a five digit octal number with a "+" or a "-" and no point preceding it. An initial "+" designates a positive number and an initial "-" a negative number; i.e. 0.12345 → +12345 and 1.12345 → -65432 respectively.

Preset Parameter:

v1: p0 to print (does not need to be inserted), or p64 to punch, or p128 to punch and print simultaneously.

Upon entering the subroutine:

AC: number to be recorded.

Temporary Registers:

d: unused
1t: register used to store the remainder of the word.
2t: register used to store the digit counter.

Enter→00	ta 17r	Set return address	(Or)17	<u>sp (0)</u>	Return to main program
01	ts 1t	Store number	2r→18	cm 1t	Positive magnitude in AC
02	<u>cp 18r</u>	Is number positive?	19	ts 1t	Store number
03	ca 23r	Put Flexo code for "+" in AC	20	ca 24r	Put Flexo code for "-" in AC
21r→04	qpl28sl	Print sign	21	<u>sp 4r</u>	
05	sr *12	Set digit Counter	22	p25r	Start of number table
06	ts 2t		23	1.30043	"+" character
16r→07	ca 1t	Value in AC	24	1.30007	"-" character
08	sr *12		25	p45	} Number Table
09	ad 22r	Add start of number table	26	p36	
10	td 13r		27	p39	
11	sl 15	Store	28	p3	
12	ts 1t	Remainder	29	p21	
(10r)13	ca (0)	Put Flexo code for digit in AC	30	p33	
14	qpl28sl	Print digit	31	p43	
15	ao 2t	Have all digits been printed?	32	p15	
16	<u>cp 7r</u>				

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Print C(AC) as Octal Number, Sign Digit and Complement, Point, Single Column Layout.			LSR OT 1.3t
			TAPE T542-3
			Classification Closed
No. of Regs. in Subroutine 36	Temp. Regs. used by Sub. 3	Average Time (operations) 8 print	Max. Time (operations) 8 print
Preset Parameters v1 P0 (does not need to be inserted) to print, or P64 to punch, or P128 to punch and print simultaneously. v2 (desired digit length) - 5			
Program Parameters on entering Subroutine ac: number to be recorded			
Results on leaving Subroutine ac: number has been recorded and is no longer available, C(AC) = qp 144 sl.			
Description When control is transferred to this subroutine, the number in the accumulator is printed as a five digit octal number preceded by either a "0" for a positive or a "1" for the complement of a negative number and an octal point. After the last digit is printed a carriage return is executed so that a single column of numbers is tabulated whenever the subroutine is used repeatedly.			
Notes: 1. The qp instructions are normally set to print and suppress the punch. By inserting p64 or p128 in preset parameter v1, one can punch and suppress the printer or punch and print simultaneously. 2. Form of output is 0.12345 appears as 0.12345 and 1.12345 appears as 1.12345.			
MSD 1/23/52	AS 2/5/52		

TITLE: PRINT C(AC) AS OCTAL NUMBER, SIGN DIGIT AND
 COMPLEMENT, POINT, SINGLE COLUMN LAYOUT

LSR: OT 1.3t
 TAPE T 542-3

Abstract: This subroutine prints a five digit octal number with a "0" or "1" and a point preceding it and then a carriage return. An initial "0" designates a positive number and an initial "1" the complement of a negative number.

Preset Parameters:

v1: p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously.
 v2: (desired digit length) - 5.

Upon entering the subroutine:

AC: number to be recorded

Temporary registers:

d: unused
 1t: register used to store the remainder of the word.
 2t: register used to store the digit counter.

00	ta 20r	Set return address	18	ca 15r	} Print a carriage return
01	ts 1t	Store number	19	qp 128s1	
02	<u>op 21r</u>	Is number positive?	(Or)20	<u>sp 0</u>	Return to main program
03	ca 27r	{ Put Flexo code for "0" and point in AC	2r → 21	ad 25r	Change sign
24r → 04	qp 134s1		Print{"0"}, shift	22	ts 1t
05	qp 128s1	Print point	23	ca 28 r	{ Put flexo code for "1" and point in AC
06	cs 35r	} Set digit counter	24	<u>sp 4r</u>	
07	ts 2t		25	0.77777	
17r → 08	ca 1t	Number in AC	26	p27r	Start of number table
09	sr* 12		27	0.07155	} "." and "0"
10	ad 26r	} Add start of number table	28	0.07144	
11	td 14r		29	p39	
12	sl 15	} Store remainder	30	p3	} Number table
13	ts 1t		31	p21	
(11r) → 14	ca (0)	{ Put Flexo code for digit in AC	32	p33	
15	qp 144s1		Print digit	33	
16	ao 2t	} Have all digits been printed?	34	p15	
17	<u>op 8r</u>		35	p4a2	

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Print C(AC) as Decimal Fraction, Sign and Magnitude, Point, Single Column Layout.			LSR OT 2.2t
			Tape T664-3
			Classification Closed
No. of Regs. in Subroutine 38	Temp. Regs. used by Sub. d - 2t (3)	Average Time (operations) 8 print	Max. Time (operations) 8 print
Preset Parameters v1 p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously v2 (Desired digit length) - 5			
Program Parameters on entering Subroutine ac: x, number to be recorded			
Results on leaving Subroutine ac: number has been recorded and is no longer available, c(ac) = +0			
Description: When control is transferred to this subroutine the number in the accumulator is printed as a five-digit decimal fraction preceded by a sign and decimal point and followed by a carriage return. The last printed digit does not contain roundoff of additional digits.			
Note: 1. The qp orders are normally set to print and suppress the punch. By inserting p64 or p128 in preset parameter v1 one can punch and suppress the printer or punch and print simultaneously.			
MD 1/18/52	JWCIII 1/21/52	AS 1/23/52	

TITLE: Print C(AC) as Decimal Fraction,
Sign and Magnitude, Point, Single
Column Layout.

LSR# OT 2.2t
Tape T664-3

Abstract: This subroutine prints a five digit decimal fraction preceded by a sign and decimal point and followed by a carriage return; i.e. 0.77776 appears as +.99993 and 1.77776 as -.00006. The last printed digit does not contain roundoff of additional digits.

Preset Parameters:

v1: p0 to print, or p64 to punch, or pl28 to print and punch simultaneously.

v2: (Desired digit length) - 5.

Upon entering the subroutine:

AC: x, number to be recorded

Temporary Registers:

d: unused

1t: register used to store remainder of the word.

2t: register used to store the digit counter.

	00	ta 20r	Set return address	(Or) 20	sp(0)	Return to main program
	01	ts 1t	Store number	2r→ 21	ca 26r	{ Put Flexo code for "-" and
	02	<u>cp 21r</u>	Is number positive?	22	<u>sp 4r</u>	dec. point in AC
	03	ca 25r	{ Put Flexo code for "+" and	23	p 10	10 x 2 ⁻¹⁵
			{ dec. point in AC			
22r→	04	qp 134 sl	Print {+}, shift	24	p 27r	Start of number table.
	05	qp 128 sl	Print decimal point	25	0.07143	+. character
	06	os 37r	} Set digit counter	26	0.07107	-. character
	07	ts 2t		27	p 45	
17r→	08	om 1t	x in AC	28	p 36	} Number table
	09	mh 23r	x · 10 x 2 ⁻¹⁵	29	p 39	
	10	ad 24r	} Add start of	30	p 3	
	11	td 14r		number table	31	
	12	sl 15	} Store remainder	32	p 33	
	13	ts 1t		33	p 43	
(11r)	14	ca(0)		{ Flexo code for	34	
			{ digit in AC			
	15	qp 128 sl	Print digit	35	p 13	
	16	ao 2t	} Have all digits	36	p 49	
	17	<u>cp 8r</u>		been printed?	37	p 4 a2
	18	ca 19r	} Print a			
	19	qp 144 sl		carriage return		

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Print Magnitude as Decimal Integer from AC, Initial Zero Suppression, No Carriage Return			LSR# OT 2.5 t
			Classification Closed
No. of Regs. in Subroutine 45	Temp. Regs. used by Sub. 4	Average Time (operations) 5 print	Max. Time (operations) 5 print
Preset Parameter v1 p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously			
Program Parameters on entering Subroutine AC: integer to be printed			
Results on leaving Subroutine Number printed out on typewriter			
<p>Description</p> <p>This subroutine prints out the magnitude of the number contained in the accumulator as a decimal integer. If $x \cdot 2^{-15}$ is the binary number in the accumulator, the decimal digits of the equivalent decimal integers d_1 are obtained as follows:</p> $d_1 = \left[(x \cdot 2^{-15}) \left(\frac{2^{11}}{10^4} \right) 2^{-11} \right] = \left[\frac{x}{10^4} \right] \cdot 2^{-15}$ <p>where $[] =$ "integral part of"</p> $d_2 = \left[\frac{x \cdot 2^{-15}}{10^4} - \left[\frac{x}{10^4} \right] 10^4 \cdot 2^{-15} \right] \left(2^4 \right) \left(\frac{10}{2^4} \right)$ $= \left[\frac{x - \left[\frac{x}{10^4} \right] \cdot 10^4}{10^3} \right] \times 2^{-15}$ <p>etc.</p> <p>This form is used so that none of the manipulations will yield an overflow. Actually, because the number $2^{11}/10^4$ is not</p>			

TITLE: Print Magnitude as Decimal Integer from
AC, Initial Zero Suppression, No
Carriage Return

LSR# OT 2.5 t

exact in the machine, a binary round-off is added at one place in the program to give a correct answer.

Notes:

1. Initial zeroes are suppressed. Zero will be printed as five spaces.
2. The form of output is :
 63
 17352
 1
 (zero)
 27
3. The program resets automatically.
4. There is no carriage return.

TITLE: Print Magnitude as Decimal Integer from LSR# OT 2.5 t
 AC, Initial Zero Suppression, No Carriage
 Return

Abstract: This subroutine prints out the magnitude of the number contained in the AC as a decimal integer. If $x \cdot 2^{-15}$ is the binary number in the AC, the decimal digits of the equivalent decimal integers d_i are obtained as follows:

$$d_1 = \left[x \cdot 2^{-15} \left(2^{11} / 10^4 \right) \left(2^{-11} \right) \right] = \left[\frac{x}{10^4} \right] \times 2^{-15}$$

where $[\]$ = "integral part of"

$$d_2 = \left[\frac{x \cdot 2^{-15}}{10^4} - \left[\frac{x}{10^4} \right] \cdot 2^{-15} \right] \left(2^4 \right) \left(\frac{10}{2^4} \right) = \left[\frac{x - \left[\frac{x}{10^4} \right] 10^4}{10^3} \right] \times 2^{-15}$$

etc.

This form is used so that none of the manipulations will yield an overflow. Actually, because the number $2^{11}/10^4$ is not exact in the machine, a binary round-off is added at one place in the program to give a correct answer.

Preset Parameters

v1 p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously

Temporary Registers

d unused
 1t temporary storage
 2t digit counter
 3t suppressor

Enter	00	ta 14r	Set return address	06	cm 21r	} $x \cdot 2^{11}/10^4$
	01	ts 21r	Store number in AC	07	mh 32r	
	02	cs 31r	} Set digit Counter	27r	08	} $x \cdot 2^{-11}$
	03	ts 2t		09	ts 21r	
	04	ca 0	} Set zero	10	sl *15	} Subtract off
	05	ts 3t		11	ts 1t	
			} Suppressor			

TITLE: Print Magnitude as Decimal Integer from LSR# OT 2.5 t
 AC, Initial Zero Suppression, No Carriage
 Return

12	ao 2t	} Advance	17r-→28	ca 29r	} Print
13	<u>_cp 15r</u>		29	qp136sl	
(Or)14	<u>sp (0)</u>	} Return to main program	30	<u>sp 25r</u>	
13r-→15	ca 21r	} Check for	31	p5	Counter
16	su 3t	} Initial	32	0.15067	$\approx 2^{11}/10^4$
17	<u>_cp 28r</u>		} Zero	33	ca 35r
18	ca 21r	} No initial zero; Add in table.	34	0.50000	10/16
19	ad 33r		} entry	35	p45
20	ts 21r	} Set	36	p36	} Table of Decimal
(1r)(20r)21	(p0)	} Printer	37	p39	
22	qp128sl	Print	38	p3	} Flexowriter Characters
23	cs 0	} Eliminate Zero	39	p21	
24	ts 3t		} Suppressor	40	p33
30r--→25	ao 1t	Round-off	41	p43	
26	mh 34r	Multiply by $1/6$	42	p15	
27	<u>sp 8r</u>	Recycle	43	p13	
			44	p49	

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Print C(AC) as Decimal Integer, Magnitude only, Initial Zero Suppression, Print Final Zero, No Layout		LSR OT 2.51t	
		TAPE T-883	
		Classification Closed	
No. of Regs. in Subroutine 49	Temp. Regs. used by Sub. 4 (d - 3t)	Average Time (operations) 5 print	Max. Time (operations) 5 print
Preset Parameters v1 p0 (need not be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously			
Program Parameters on entering Subroutine AC: Integer to be printed			
Results on leaving Subroutine Number printed out on typewriter			
Notes: <ol style="list-style-type: none"> 1. Initial zeroes are suppressed. Zero will be printed out as four spaces and a zero. 2. The form of output is: <div style="text-align: center; margin: 10px 0;"> 63 17352 1 ...->>>>0 (zero) 27 </div> 3. The program resets automatically. 4. There is no carriage return. 			
JWC III 2/10/52	MAS 2/11/52		

TITLE: Print C(AC) as Decimal Integer, Magnitude only,
Initial Zero Suppression, Print Final Zero, No
Layout

LSR OT 2.51t
TAPE T-883

Abstract: This subroutine prints out the magnitude of the number contained in the accumulator as a decimal integer. If $x \cdot 2^{-15}$ is the binary number in the accumulator, the decimal digits of the equivalent decimal integers d_i are obtained as follows:

$$d_1 = \left[\left(x \cdot 2^{-15} \right) \left(\frac{2^{11}}{10^4} \right) \left(2^{-11} \right) \right] = \left[\frac{x}{10^4} \right] \cdot 2^{-15}$$

where $[] =$ "integral part of"

$$d_2 = \left[\left(\frac{x \cdot 2^{-15}}{10^4} - \left[\frac{x}{10^4} \right] \cdot 10^4 \cdot 2^{-15} \right) \left(\frac{2^4}{2^4} \right) \left(\frac{10}{2^4} \right) \right]$$

$$= \left[\frac{x - \left[\frac{x}{10^4} \right] \cdot 10^4}{10^3} \right] \cdot 2^{-15}$$

This form is used so that none of the manipulations will yield an overflow. Actually, because the number $2^{11}/10^4$ is not exact in the machine, a binary round-off is added at one place in the program to give a correct answer.

Temporary Registers

d - unused
1t - temporary storage
2t - digit counter
3t - suppressor

Preset Parameters

v1 p0 (need not be inserted) to print, p64 to punch, or pl28 to print and punch

00	ta 19r	Set return address	14	ca 25r	Last digit
01	ts 25r	Set no. in AC	15	ad 37r	Add in start of table
02	cs 35r	} Set digit counter	16	ts 17r	
03	ts 2t		(16r) 17 (p0)	} Print	
04	ca 0	} Set zero suppressor	18		qp 128sl
05	ts 3t		(Or) 19	sp ()	Return to main program
06	cm 25r	} $x \cdot \frac{2^{11}}{10^4}$	13r → 20	ca 25r	} Not at last digit
07	mh 36r		21	su 3t	
31r → 08	sr* 11	} $x \cdot 2^{-11}$	22	cp 32r	
09	ts 25r		23	ad 37r	} Add in table entry
10	sl* 15	} Subtract off integral part	24	ts 25r	
11	ts 1t		(24r) 25 (p0)	Set Printer	
12	ao 2t	} Advance counter	26	qp 128sl	Print
13	cp 20r	} Are we at last digit?	27	cs 0	} Eliminate

TITLE: Print C(AC) as Decimal Integer, Magnitude only,
 Initial Zero Suppression, Print Final Zero, No
 Layout

LSR OT 2.51t
 TAPE T-883

	28	ts 3t	zero suppressor	39	p45	}	Table of Digits
34r →	29	ao 1t	Round off	40	p36		
	30	mh 38r		41	p39		
	31	<u>sp 8r</u>		42	p3		
22r →	32	ca 33r	} Print space	43	p21		
	33	qp 136sl		44	p33		
	34	<u>sp 29r</u>		45	p43		
	35	p4		46	p15		
	36	0.15067	$\frac{2.11}{10^4}$	47	p13		
	37	//ca 39r	10^4	48	p49		
	38	0.50000	$\frac{10}{16}$				

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Print Out Flexowriter Characters, Previously Stored Three Characters to a Register by LSR OT 3.lat, Using Program Parameter		LSR OT 3.lt	
		TAPE T782-2	
		Classification Closed	
No. of Regs. in Subroutine 38	Temp. Regs. used by Sub. d = 4t (5)	Average Time (operations) 13 + 14n (n=no. of chars.) (Actually limited by number of printouts)	Max. Time (operations) 13+20n (n=no. of chars.) (Actually limited by no. of printouts)
Preset Parameters v1: p0 (does not need to be inserted) to print, p64 to punch, and p128 to print and punch simultaneously.			
Program Parameters on entering Subroutine u1: p(address of first register in storage in which characters have been previously stored.)			
Results on leaving Subroutine Flexowriter characters printed out from storage.			
Description This subroutine prints out Flexowriter characters previously stored by LSR OT 3.lat, three to a register. Subroutine OT 3.lat stores characters 5-digits long, with a 00000 "tag" character indicating when the final digit is to be a 0 and when a 1. (For further description of LSR OT 3.lat, see the Specification sheet for LSR OT 3.lat). This subroutine reverses the process, unstacking and printing out. It will print out characters, placing a zero in the sixth digit, until a "tag" character is reached. Then it will place 1's in the sixth digit until another "tag" character is reached, etc. This is the exact reverse of LSR OT 3.lat. When a negative number is reached (previously stored by the read-in routine) the routine stops printing and returns control to the main program.			
Note For notes on how many characters may be stored, etc., see the Specification Sheet for LSR OT 3.lat.			

JWC III 2/1/52	AS 2/5/52
-------------------	--------------

TITLE: Print Out Flexowriter Characters,
Previously Stored Three Characters
to a Register by LSR OT 3.1t, Using
Program Parameter

LSR# OT 3.1t
TAPE# T782-2

Upon entering the subroutine:

sp (subroutine)
p (first address to be printed out of)

Preset Parameters:

v1: p0 (does not need to be inserted) to print, or p64 to
punch, or pl28 to print and punch simultaneously

Temporary Registers:

d: temp. character storage
1t: 3 counter
2t: operational register
3t: find digit
4t: digit to be exchanged

	00	ta 1r	Plant ul address		19	sl* 15	} Shift off left and restore reg. contents
(3r)(0r)	01	ca (0)	} Plant first address		20	ts 2t	
	02	td 11r				21	ca d
	03	ao 1r	} Set return address		22	sl* 1	
	04	td 13r				23	ad 3t
	05	ca 0	} Set final digit		24	qp 128sl	Print
	06	ts 3t			35r →	25	ao 1t
	07	ca 36r	} Set digit to be exchanged to pl		26	<u>cp 14r</u>	
	08	ts 4t				27	ao 11r
28r →	09	cs 37r	} Set 3-counter		28	<u>sp 9r</u>	
	10	ts 1t			17r →	29	ca 2t
(2r)	11	ca (0)	} Pick up register to be printed		30	sl* 5	
	12	ts 2t				31	ts 2t
(4r)	13	<u>cp (0)</u>	Are we finished?		32	ca 4t	} Exchange final digit
26r →	14	ca 2t	} No. Shift character into position		33	ex 3t	
	15	sr* 10				34	ts 4t
	16	su 0	} Is this character a zero?		35	<u>sp 25r</u>	} Back to check finish final digit
	17	<u>cp 29r</u>				36	
	18	ts d	No		37	p2	character counter

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Read in Standard Flexwriter Characters and Store Three Characters to a Register, Using Program Parameter. (For Use with LSR OT 3.lt.)			LSR OT 3.1at
			TAPE T781-2
			Classification Closed
No. of Regs. in Subroutine 66	Temp. Regs. used by Sub. $d - 6t$ (7)	Average Time (operations) $40 + 13n$, where n is no. of characters (Present limiting case is no. of read-ins on PETR.)	Max. Time (operations) $40 + 26n$ (in the most unusual case) (Present limiting case is no. of read-ins on PETR)
Program Parameters: on entering Subroutine ar: u_1 (automatic) ul p (address of first register in which characters are to be stored)			
Results: on leaving Subroutine Flexwriter characters stored, five binary digits to the character, beginning at the register given in u_1 , three characters to the register. All registers are positive except the one succeeding the final storage register, which is set to $n0$.			
Description This subroutine reads in Flexwriter characters continuously from the photoelectric reader, and stores them in the registers beginning at the address stored in u_1 . The standard Flexwriter tape to be called in is typed on a Flexowriter having a binary code the same as the code on the Flexowriter connected to the Whirlwind output. A "stop" character, <u>with a seventh hole</u> , must be typed at the end of the tape as a "tag" character so that the subroutine may stop the storage process.			
Notes 1. The Flexwriter characters are shortened to five-digit length, and stored three to a register. A second "tag" character, 00000, is used to indicate whether or not the sixth digit in a character is to be a zero or one. Whenever a series of characters switches from final digit "zero", to final digit "one", and vice versa, such a tag is inserted. This is not necessarily a waste of storage, since all alphabetical characters in the standard Flexwriter code have the sixth hole zero, and thus there is no need for such a tag in most storage of straight language. Similarly, all numbers and most punctuation marks have the sixth hole a "one", and so a series of numbers does not require such a "tag" character until an alphabetical character appears. Mixtures of constantly changing letters and			

TITLE: Read in Standard Flexowriter Characters and Store
Three Characters to a Register, Using Program
Parameter. (For Use with LSR OT 3.lt.)

LSR OT 3.lat
TAPE T781-2

numbers will, however, be wasteful of storage; and in such rare cases, another pair of storage and print routines should be used. With ordinary typescript, an estimated 25% of storage can be gained, however, over "two-to-a-register" methods of storing Flexowriter standard characters.

2. The number of registers needed to store a given tape will thus be a function of the number of changes from "sixth-hole-zero" to "sixth-hole-a-one". In standard English language, an average formula would seem to be

$$N = \frac{n}{3} + \frac{n}{20} = \frac{23n}{60} = .38n$$

Where N is the number of registers needed for storage, and n the number of characters, $n/20$ being considered an average number of such changeovers.

3. This subroutine must always be used in conjunction with Print-Out Routine OT 3.lt. It may be used in the following manner: Read in subroutine OT 3.lt, preferably in the upper or lower ends of storage. When control is transferred to this subroutine it will automatically read in and store the required standard Flexowriter characters in the addresses beginning at the address given by program parameter ul. After the final stop character is reached, it will not store this, but instead make the succeeding register a negative zero. Control is then automatically transferred to the instruction following the ul register.

Following this, OT 3.lt may be read in by means of the standard read-in conversion program, or else, if it has been previously stored, control may be transferred to it at any desired time. For a fuller explanation of how this program prints-out, see the Specification Sheet for OT 3.lt.

4. The character 000001 cannot be read in by this subroutine, since it will be confused by the machine with the 5-digit "tag" character 00000. However, in the 1951 Flexowriter code, this is "back-space", which should not be needed, and in the 1952 Flexowriter code, this character is not produced by any key on the keyboard.

JWC III 1/24/52	AS 2/5/52
--------------------	--------------

TITLE: Read in Standard Flexowriter Characters
and Store Three Characters to a Register,
Using Program Parameter. (For Use with
LSR # OT 3.1t.)

LSR# OT 3.1at
TAPE# T781-2

Upon entering the subroutine:

sp: subroutine
p: first address to be stored

Temporary registers:

d: temporary digit store
lt: last 6th digit representation
2t: character store
3t: present 6th digit representation
4t: temporary store
5t: 3-counter
6t: temporary store

00	ta 1r	Plant ul address	33r →	20	ca 2t	} Store present 6th digit representation
(3r)(0r)01	ca ()	Plant 1st address where machine is to store chars.	21	sl* 14	} Are past and present 6th digits the same?	
02	td 53r		22	sr* 14		
03	ao 1r	Set return address	23	su 0		
04	td 42r		24	ts 3t		
05	ca 0	Set last digit reg. in for zero	25	mh 1t	} Store 5 digits of character	
06	ts 1t		26	<u>op 43r</u>		
07	ca 64r	Store 3 counters	47r →	27	ca 2t	} Recycle
08	ts 5t		28	<u>sp 48r</u>		
09	ca 0	Set zero in temp store	29	<u>sp 11r</u>	} Is char. a "stop"?	
10	ts 4t		19r →	30		ca 65r
29r →	11	qr 0	Read in	31	su 2t	} No. Back to check digits
12	ca 8	} Store 6 digits in 2t	32	<u>op 34r</u>		
13	sr* 5		33	<u>sp 20r</u>		
14	td d		32r →	34	ca 0	} Yes.
15	ca d		35	<u>sp 48r</u>	Store 00000	
16	sr* 5	} Store 00000	36	ca 0		
17	ts 2t		37	<u>sp 48r</u>		
18	su 65r	} Check to see if char. is a "stop"	38	ca 53r	} Plant last register address	
19	<u>op 30r</u>		39	td 41r		

TITLE: Read in Standard Flexowriter Characters
and Store Three Characters to a Register,
Using Program Parameter. (For Use with
LSR # OT 3.1t.)

LSR# OT 3.1at
TAPE# T781-2

	40	cs 0	} Put negative no. in last reg.	56	ao 5t	} Are we at end of register?
39r →	41	ts (0)		57	<u>cp 63r</u>	
(04r) →	42	<u>sp (0)</u>	Back to main Program	58	ca 0	} Yes. Reset temp. store.
26r →	43	oa 3t	} Switch digit representations	59	ts 4t	
	44	ts 1t		60	cs 64r	Reset 3-counter
	45	ca 0	} Store 00000	61	ts 5t	
	46	<u>sp 48r</u>		57r ↓	62	ao 53r
	47	<u>sp 27r</u>	Back to store 5 digits (43r) 57r ↓	63	<u>sp (0)</u>	Exit from internal s.r.
28r, 37r } 35r, 46r }	48	ta 63r	Plant link of internal subroutine	64	p2	3-counter reset
	49	sr* 1	} Shift left and store	65	p5l	"stop" character
	50	ts 6t				
	51	ca 6t				
	52	ad 4t				
(62r)(02r)	53	ts (0)				
	54	sl* 5				
	55	ts 4t				

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: (30,0,0) MRA Print and/or Punch, Decimal Fraction, Sign, Number of Digits Arbitrary, No Carriage Return, Sign Agreement (Interpreted)			LSR OT 103.10t
			TAPE T799-1
			Classification Closed, Interpreted
No. of Regs. in Subroutine 76	Temp. Regs. used by Sub. 4	Average Time (operations) 94+19n, number of digits printed = n	Max. Time (operations) 94 + 27n
<p>Preset Parameters</p> <p>vx: pN, where N is the address in storage of the first register of the interpretive subroutine (in title of main program)</p> <p>vl: p0 (does not need to be inserted) to print, p64 to punch, or pl28 to print and punch simultaneously</p> <p>v2: pn, where n is the number of decimal digits to be printed.</p>			
<p>Description</p> <p>This subroutine prints and/or punches the sign and magnitude of the contents of the MRA in the following manner</p> $\begin{matrix} + \\ -.d_1d_2 \dots d_n \end{matrix}$ <p>The number, n, of decimal digits to be printed is a preset parameter (v2). The digits, d_i, are obtained by multiplying the magnitude of the contents of the MRA successively by pl0.</p> <p>This subroutine contains a sign agreement program so that the contents of the MRA need not be a number whose major and minor parts are of like sign.</p> <p>The sp instruction transferring control to this subroutine must be an interpreted sp (i.e., control must be in the interpretive subroutine). After execution of the subroutine control remains in the interpretive subroutine which then proceeds to interpret the instruction following the sp instruction in storage.</p> <p style="text-align: center;">There is no carriage return.</p> <p>This subroutine can be used with any(30,0,0) interpretive subroutine. The contents of the MRA are left undisturbed during the execution of this subroutine.</p>			
FCH 2/1/52	MRS 2/8/52		

TITLE: MRA Print and/or Punch in (30,0,0) Interpretive LSR OT 103.10t
 Subroutine without Sign Agreement, Decimal Number, TAPE T799-1
 Sign, Number of Digits Arbitrary, No Carriage
 Return (Interpreted)

Abstract: This subroutine prints out a ± sign and a decimal point followed by the magnitude of the contents of the MRA as a decimal fraction. The decimal digits are obtained by multiplying the contents of the MRA successively by p10. The number of digits to be printed is a preset parameter (v2). There is no carriage return. The subroutine is interpreted and can be used with any (30,0,0) interpretive subroutine.

Preset Parameters

- vx: pN, where N is the address in storage of the first register of the interpretive subroutine (in title of main program)
- vl: p0 (does not need to be inserted) to print, or p64 to punch, or p128 to punch and print simultaneously
- v2: pn, where n is the number of decimal digits to be printed

Temporary Storage

- d unused
- 1t } Temporary storage
- 2t }
- 3t }
- 4t Digit counter

00	ta 37r		.35r → 17	cm 1t	} Multiply C(1t,2t) by p10
01	sp ax	Resume ordinary mm operation	18	mh 62r	
02	ca 3ax	} Store C(mra) in 1t and 2t	19	ts 3t	
03	ts 2t		20	sl 15	
04	ca 2ax		21	ts 1t	
58r → 05	ts 1t		22	cm 2t	
54r → 06	mr 2t	Do 1t and 2t agree in sign?	23	mh 62r	
07	cp 38r		24	ts 2t	
08	ca 1t	} Sense and print algebraic sign of C(mra) followed by a decimal point.	25	sl 15	
56r → 09	cp 12r		26	ex 2t	
10	ca 74r		27	sa 1t	
11	sp 13r		28	ts 1t	
9r → 12	ca 75r		29	ca 3t	
1r → 13	qp 134s1	} Set up entry into table	30	ad 63r	
14	qp 128s1		31	td 32r	
15	cs 61r	} Set up digit counter	(31r) 32	ca (0)	
16	ts 4t		33	qp 128s1	

} Print a single digit

TITLE: MRA Print and/or Punch in (30,0,0) Interpretive Subroutine without Sign Agreement, Decimal Number, Sign, Number of Digits Arbitrary, No Carriage Return (Interpreted)

LSR OT 103.10t
TAPE T799-1

	34	ao	4t	} Have enough digits been printed?	48r →	55	ao	1t		
	35	<u>cp</u>	<u>17r</u>				56	<u>sp</u>	9r	
	36	<u>sp</u>	<u>ax</u>	} Return control to int. subroutine	40r →	57	cs	1t		
(Or)	37	<u> sp</u>	<u>(0)</u>				58	<u>sp</u>	5r	
7r →	38	cm	1t	} Is C(1t) ≠ 0?		59	<u>ca0</u>			
	39	su	0				60	<u>pl</u>		
	40	<u>cp</u>	<u>57r</u>				61	<u>nla2</u>	v2	
	41	cm	2t	} Is C(2t) ≠ 0?		62	<u>pl0</u>			
	42	su	0				63	<u>p64r</u>		
	43	<u>cp</u>	<u>51r</u>				64	<u>p45</u>	0	Table
	44	su	60r	} Form 1 - C(2t)		65	<u>p36</u>	1		
	45	<u>ad</u>	<u>59r</u>				66	<u>p39</u>	2	
	46	ts	2t				67	<u>p3</u>	3	
	47	ca	1t	} Is C(1t) pos.?		68	<u>p21</u>	4		
	48	<u>cp</u>	<u>55r</u>				69	<u>p33</u>	5	
	49	su	60r	} Form C(1t) - 2 ⁻¹⁵		70	<u>p43</u>	6		
	50	ts	1t				71	<u>pl5</u>	7	
43r →	51	cs	2t	} Complement C(2t)		72	<u>pl3</u>	8		
	52	ts	2t				73	<u>p49</u>	9	
	53	ca	1t	} Re-enter sign agree- ment		74	0.07143	+		
	54	<u>sp</u>	<u>6r</u>				75	0.07107	-	

PROGRAMMED ARITHMETIC

PA 1.2 t (T747-3) is the same as PA 1.1 t except that
register 26r reads pax2.

PA 2.2 t (T723-1) is the same as PA 2.1 t except that
register 197r reads pax2.

Both the above subroutines use the new Flexowriter code
whereas PA 1.1 t and PA 2.1 t both use the old
Flexowriter code.

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on Real (15,15,0) Floating Point Double Register Numbers (General Subroutine)			LSR PA 1.1 t
			Classification Interpretive
No. of Regs. in Subroutine 97	Temp. Regs. used by Sub. -----	Average Time (operations) see description	Max. Time (operations) see description
Preset Parameters (to be typed in the title of the main program) v8/ pk: k = separation between two registers of the number vx/ pN: N = location of first register of subroutine in storage			
<p>Description:</p> <p>By means of this subroutine various logical and arithmetic operations can be carried out upon real numbers represented in the (15,15,0) system and stored in two registers, say <u>n</u> and <u>n+k</u>, where <u>k</u> is specified by a preset parameter.</p> <p>If <u>a</u> is any number such that $2^{-15} < a < 2^{15}$, or $a = 0$, then there exists a signed 15 binary digit number <u>x</u> and a signed 15 binary digit integer <u>y</u> such that</p> $\frac{1}{2} < x < 1 \text{ or } x = 0,$ $0 \leq y < 2^{15},$ <p>and $\left \frac{a - x \cdot 2^y}{a} \right \leq 2^{-16}.$</p> <p><u>x</u> and <u>y</u> are stored respectively in the two registers <u>n</u> and <u>n+k</u> for use by this subroutine.</p> <p>Example (a) Let $a = -300$ (decimal) Then $a = -100101100$ (binary) and hence $x = 1.011010011111111$ and $y = 1001$ Register <u>n</u> contains 1.011010011111111 and Register <u>n+k</u> contains 0.000000000001001</p>			

TITLE: Operations on Real (15,15,0) Floating
Point Double Register Numbers (General
Subroutine)

LSR PA 1.1 t

Example (b) Let $a = \frac{1}{10}$ (decimal)

Then $a = 0.000110011001100110011\dots$
and hence $x = 0.110011001100110$,
and $y = -11$

Register n contains 0.110011001100110
and Register $n+k$ contains 1.111111111111100

The programmer need not carry out this conversion process himself, but instead need only find a five decimal digit X and Y such that

$$\frac{1}{10} \leq X < 1, \text{ or } X = 0$$

$$0 \leq |Y| \leq 32,767$$

$$\text{and } \left| \frac{a - X \cdot 10^Y}{a} \right| \leq 10^{-5}$$

By giving this information to the proper subroutine in the IP section of the Library of Subroutines, the remainder of the conversion can be done automatically.

Operations upon numbers in this representation are coded using an instruction code similar to but differing somewhat from the standard WW code. Any number of these instructions can be performed in sequence by placing an sp ax before the first instruction of the sequence. The instructions are then interpreted successively until a change-of-control instruction is reached at which point another sequence is performed unless the instruction is an sp ax. In the latter case, ordinary WW operation is resumed at the register following the sp ax.

The subroutine contains the analogy of an accumulator and a program counter. The "program counter" is in register 76r and contains the address of the register containing the instruction being interpreted. The "multiple register accumulator" (mra) is in registers 2r, and 3r and contains the result of the last instruction (except ta, cp, and sp). Register 2r contains x and 3r contains y .

TITLE: Operations on Real (15,15,0) Floating
Point Double Register Numbers (General
Subroutine)

LSR PA 1.1 t

Instruction

Function

ts n	Store the number in the mra in registers <u>n</u> and <u>n+k</u> .
ta n	Store the address of the register following the last sp or cp instruction in the address section of register <u>n</u> .
ex n	Exchange the number in the mra with the number in registers <u>n</u> and <u>n+k</u> .
cp n	If the number in the mra is negative, proceed as in sp, otherwise ignore the instruction.
sp n	If $n \neq ax$, interpret next the instruction in register <u>n</u> . If $n = ax$, resume ordinary WW operation at the register following the register containing the <u>sp ax</u> .
ca n	Clear the mra and then add into the mra the number in registers <u>n</u> and <u>n+k</u> .
cs n	Clear the mra and then add the negative of the number in registers <u>n</u> and <u>n+k</u> into the mra.
ad n	Add the number in registers <u>n</u> and <u>n+k</u> to the number in the mra and leave the result in the mra.
su n	Subtract the number in registers <u>n</u> and <u>n+k</u> from the number in the mra and leave the result in the mra.
cm n	Clear the mra and add the magnitude of the number in registers <u>n</u> and <u>n+k</u> into the mra.
mr n	Multiply the number in the mra by the number in registers <u>n</u> and <u>n+k</u> and leave the result in the mra.
dv n	Divide the number in the mra by the number in registers <u>n</u> and <u>n+k</u> and leave the result in the mra.

N.B. In the above the phrase "the number in" should be read to mean "the number represented by the contents of".

TITLE: Operations on Real (15,15,0) Floating
Point Double Register Numbers (General
Subroutine)

LSR PA 1.1 t

Notes:

1. Entering and Leaving the Subroutine

Wherever an sp ax is encountered, the machine will enter the subroutine if it is not already in it, or leave it if it is in the subroutine.

2. Preset Parameters

A general preset parameter, vx, is used to specify the location in storage of the first register of the subroutine. Another general preset parameter, v8, is used to specify the separation k of the two registers containing a number. Both of these parameters must be specified in the title of the program tape.

3. Accuracy of Arithmetic Operations

The operations multiply and divide give a result with at least $14 \frac{1}{2}$ digit accuracy, i.e. the result is rounded off at worst in the 15th place. The relative error in the operations add and subtract is at worst 2^{-15} , but it may be biased downwards if the magnitude of the sum is sufficiently greater than the magnitude of either of the addends.

4. Alarms

Arithmetic overflow alarms can occur in registers 19r, 20r, 47r, 71r, or 72r. In all cases this occurs when the exponent y of the result being calculated during an ad, su, mr, or dv instruction lies outside the range $(-2^{15}, 2^{15})$. When such an alarm occurs, the address of the instruction that is being interpreted will be in register 76r. Any other alarms in the subroutine will be due to excessive addresses or use of an instruction not listed above.

TITLE: Operations on Real (15,15,0) Floating Point Double
Register Numbers (General Subroutine)

LSM PA 1.1 t

Instruction Code and Operation Times:

ts 23	cp 17(+), 27 (-)	cs 26	cm 25
ta 18	sp 24	ad 41	mr 23
ex 23	ca 27	su 43	dv 26

Reset Parameters: (to be typed in the title of the main program)

v8/ pk: k = separation in storage of two registers of numbers
vx/ pN: N = address in storage of initial register of subroutine

00	ta 76r	} Store address of 1st instruction to be interpreted	25	(p0)	digit storage
01	<u>sp 76r</u>		26	pa8	separation parameter
02	(p0)	} multiple register accumulator	89r-→27	<u>sp 90r</u>	"ts"
03	(p0)		28	sp 19r	
89r-→04	ca ax	"ca"	89r-→29	ca 42r	} Store digits in A-register
89r-→05	<u>sp 8r</u>	"cs"	(78r)30	ta (0)	
89r-→06	<u>sp 46r</u>	"ad"	31	<u>sp 75r</u>	
89r-→07	<u>sp 44r</u>	"su"	89r-→32	<u>sp 90r</u>	"ex"
5r,89r-→08	ca 86r	"cm" Pick up and store	89r-→33	<u>sp 22r</u>	"cp"
09	ts 93r	} $\left. \begin{matrix} cs \\ ca \\ cm \end{matrix} \right\} n+k$	24r,89r-→34	ao 76r	} set A-register address
10	<u>sp 90r</u>		35	td 42r	
11	pl6		36	ca 88r	} set pick-up instruction's address
89r-→12	mr 2r	"mr" Form $x_1 \cdot x_2$	37	td 76r	
13	<u>sp 69r</u>		38	su 4r	
89r-→14	ex 2r	} Form and Store $\frac{x_1 \cdot 2^{-1}}{x_2} \cdot 2$ s.f.	39	ts 95r	} Does address of cp or sp instruction differ from <u>ax</u> ?
15	sr 1		40	cm 95r	
16	dv 2r		41	su 0	
17	sf 25r		(35r) 42	<u>cp (0)</u>	
18	ex 3r		43	<u>sp 76r</u>	
19	su 95r	} Form $y_1 - y_2 + 16$	7r-→44	ts 91r	} Complement x_2
20	ad 11r		45	cs 91r	
21	<u>sp 72r</u>	} Is $x_1 < 0$?	6r-→46	ex 95r	} Form and store $y_2 - y_1$.
33r-→22	cs 2r		47	su 3r	
23	<u>cp 75r</u>		48	ts 91r	
24	<u>sp 34r</u>	Go to sp routine	49	<u>cp 55r</u>	Is $y_2 - y_1 > 0$?

TITLE: Operations on Real (15,15,0) Floating Point Double LSR# PA 1.1 +
 Register Numbers (General Subroutine)

50	ad 3r	} Interchange (x_1, y_1) and (x_2, y_2)	94r→73	ex 3r	} Store result in proper registers Increase address of pick-up instruction Pick-up instruction
51	ts 3r		74	ts 2r	
52	ca 2r		23r,31r,5r→75	ao 76r	
53	ex 95r		(37r)(0r)1r,3r→76	ca (0)	
54	ts 2r		77	ts 91r	} Store $ y_2 - y_1 $ and digits where necessary
49r→55	cm 91r		78	td 30r	
56	td 62r		79	td 88r	
57	su 96r	Is $ y_2 - y_1 > 15$?	80	ad 26r	
58	<u>cp 60r</u>		81	ts 93r	} Form <u>sp</u> to proper part of subroutine
59	<u>sp 75r</u>	Addition unnecessary	82	td 86r	
58r→60	ca 96r	Set $y_2 = 15$	83	sr *11	
61	ex 95r		84	ad 28r	
(56r)62	sr (0)	Form $x_2 2^{-y_2 - y_1}$	85	td 89r	} Pick up and store y_2 Pick up x_2 Go to part of subroutine for particular instruction
63	sa 2r	} Add	(82r)86	ca (0)	
64	ts 2r		87	ts 95r	
65	ca 0	} Add overflow into 2^{-15} times the sum	(79r)88	ca (0)	
66	ex 2r		(85r)89	<u>sp (0)</u>	
67	sr *15		10r,32r→90	ca 2r	
68	ad 2r		(77r)91	(p0)	
13r→69	sf 25r	Scale factor	92	ex 3r	} 4r-ca 14r-dv 5r-cs 27r-ts 6r-ad 29r-ta 7r-su 32r-ex 8r-cm 33r-cp 12r-mr 34r-sp
70	ex 3r	and store result (9r), (81r)93	(p0)		
71	ac 96r	Form $y_2 + y_1$	94	<u>sp 73r</u>	
21r→72	su 25r	Subtract scale factor	95	(p0)	
			96	p15	Temporary storage

DIGITAL COMPUTER LABORATORY

MASSACHUSETTS INSTITUTE OF TECHNOLOGY

Specifications of WHIRLWIND I LIBRARY SUBROUTINE Number PA 2.2

Title: Extra-Precision and Floating-Point Real Number Arithmetic, using 2-register 24,6,0 Numbers; Basic Instruction Code with Division, INTERPRETIVE

Total Number of Registers Occupied by the Subroutine: 204 storage registers
 Temporary Storage Registers Required by the Subroutine: no temporary regs.
 Time Required to Perform the Subroutine: average = 50* WNI operations
 maximum = 76* WNI operations

* per interpreted operation; see page 4 for details

Preset Parameters (Values to be indicated in tape title line)

x | pN: N = address assigned to the initial register of the subroutine
 x2 | pk: k = separation between registers assigned to each 2-register number

Description

This interpretive subroutine, when called into action, takes instructions (more strictly, program parameters written as instructions) one at a time from consecutive storage registers and performs the designated single-address operations defined by the interpreted-instruction code given on page 4. These operations are primarily arithmetical operations performed on real numbers represented in the 24,6,0 system. Each number is stored in some multiple-register location n consisting of the pair of registers n and n+k, where n is the address of the given location and k is determined by preset parameter x2.

The 24,6,0 number system represents any real number N, provided that either $N=0$ or $2^{63} > |N| \geq 2^{-64}$, as a signed 24-binary-digit fraction x and a signed 6-binary-digit integer y, where x and y are chosen in such a way that either $x=0$ or $1 > |x| \geq .5$ and that $|1 - x2^y N^{-1}| < 2^{-24}$. Thus the number pair x,y represents N to within $\pm 0.000006\%$, equivalent to about 7 significant decimal digits. The sign and first 15 digits of x occupy one register while the sign and 6 digits of y and the last 9 digits of x occupy the second register of the pair assigned to contain the number N. Details of this and other number systems are available elsewhere.

A multiple-register accumulator (MRA) is used in place of the AC in many interpreted operations. This MRA is not a special register as is the AC but rather is a group of 3 ordinary storage registers contained within the interpretive subroutine, specifically registers 2r, 3r, and 4r. Even though only 2 registers are needed to contain a 24,6,0 number, 3 registers are used for the MRA to avoid the time-consuming operation of packing the last 9 digits of the number and the sign and 6 digits of the exponent together into one register after each interpreted instruction. A further advantage is gained in that any sequence of arithmetic operations is performed using 30 digits for the number and 15 digits for the exponent. This provides in effect a 30,15,0 system. The 24 and 6 limitation is imposed only when necessary, namely on ts and ex operations. Thus greater range and greater precision are available in sequences of arithmetic operations than the 24,6,0 system would normally allow.

The roundoff error on ad and su is made in the 29th digit of the sum before it is scale-factored. That is, in adding any two 24,6,0 numbers, $v \cdot 2^w$ to $x \cdot 2^y$, assuming $1 > |v| \geq .5$, $1 > |x| \geq .5$, $w \geq y$, the sum obtained is $[(v + x \cdot 2^{y-w} + 2^{-29}) 2^z] 2^{w-z} = u \cdot 2^{w-z}$, where z is chosen in such a way that $1 > |u| \geq .5$.

The roundoff in mr is made in the 28th digit.

The roundoff in dv is made in the 27th digit.

The roundoff in ts and ex (i.e. in packing the 30,15,0 numbers into 24,6,0 form) is of course made in the 25th digit. If the exponent y is less than -63, the value -63 is substituted for it, without changing x in any way.

Arithmetic alarms, because of the floating point system employed, and because of the extended range allowed within the MRA, will normally not occur in an interpreted program unless the contents of the MRA, call it $x \cdot 2^y$, prior to a ts or ex operation has an exponent $y > 63$, in which case an overflow alarm always occurs at register 203r during the interpretation of the ts or ex operation, even if $x = 0$. If during an arithmetic operation the exponent y exceeds the bounds $2^{15} > y > 2^{-15}$, an overflow alarm will occur at register 28r, 85r, 130r, 175r or 176r.

Entry to and exit from the subroutine is accomplished by means of the instruction spax. The first instruction in a program is always performed in the Whirlwind code. When 24,6,0 operations are needed, control is transferred to the subroutine by spax, x being the parameter which specifies the location of the subroutine. Instructions following the first spax are then performed in the interpreted code. When operations on 1-register fixed-point words are desired, control is transferred back to the main program by spax. This spax is given a special interpretation by the subroutine and results in the instructions following it being performed in the Whirlwind code. Use of a sequence of Whirlwind-coded instructions between two interpreted instructions does not affect the contents of the MRA, but use of any interpreted instruction does affect the contents of the AC.

For numerical input at the present time, all decimal numbers to be converted to 24,6,0 form must be written as a signed decimal fraction which is less than 1.0 and not less than 0.1 followed by a single signed decimal digit indicating the actual position of the decimal point. That is, any number N is written in the form $N = X \cdot 10^Y$, with $1 > |x| \geq .1$ and $-9 \leq Y \leq 9$, and with X having at most 8 decimal digits. For example,

the number 300, which equals $.3 \times 10^3$, is written as $+.3 | +3$;
 the number $.01\pi$, which equals $.31415927 \times 10^{-1}$ is written as $+.31415927 | -1$;
 the number $-1/128$, which equals $-.78125 \times 10^{-2}$ is written as $-.78125 | -2$

Alternatively, any number may be converted to 24,6,0 binary form by hand and written as 2 standard single length octal numbers. The procedure for converting by hand is described elsewhere.

Allocation of storage locations to the necessary 2-register numbers, (both for the main program and the subroutines), temporary storage, the main program, the subroutines, and the interpretive subroutine PA 2.2 must at present follow a rather inflexible rule because of the input conversion procedures currently in use. The scheme to be followed is shown diagrammatically below, with decimal addresses used throughout. Notice that parameter x is at present assigned the value 852 in all programs.

<u>Numbers designated by programmer</u>	<u>Storage registers</u>	<u>Comments</u>
address at start of program, usually 32.	main program 2-register numbers, 1st halves	the assignments to consecutive locations of the 2-register constants needed by individual subroutines is handled automatically by the conversion program. The number of locations needed is the sum of the numbers needed by individual subroutines.
total number of locations = $k =$ parameter x^2 .	subroutine 2-register numbers, 1st halves	
address of start of temporary storage = parameter 0.	temporary storage, 1st halves	the number of temporary locations needed is the maximum of the numbers needed by the main program and the subroutines. Note that all locations are 2-register locations. For 1-register temporary storage, both halves of any 2-register location n may be used by referring to nt for the first half and to $ntax2$ for the second half.
	main program 2-register numbers, 2nd halves	
	subroutine 2-register numbers, 2nd halves	address of 2nd half of last main program number must be less than 530.
	temporary storage, 2nd halves	
	main program	
address of start of main program and of each subroutine and address of first instruction to be performed must be indicated	subroutines	address of last word of last subroutine must be less than 704.
		space available for print subroutine OT 102.1
address of start of interpretive subroutine = 852 = parameter x	interpretive subroutine	

The interpreted instruction code of this subroutine is given below. The instructions have the same binary value as the similar Whirlwind instructions. Hence they are written, typed and converted in the same way as Whirlwind instructions and are in fact indistinguishable from them. The term "number in location n" is used to signify the number represented in 24,6,0 form by the 32 binary digits contained in the pair of registers n and n+k. The term "register m" is used to signify the single register m. Figures in parentheses give the number of Whirlwind instructions required to interpret the indicated instructions.

<u>Interpreted Instructions</u>	<u>Function</u>
ca n (38)	Clear the MRA and add into it the number in location n.
cs n (36)	Clear the MRA and subtract from it the number in location n.
cm n (37)	Clear the MRA and add into it the magnitude of the number in location n.
ad n (72)	Add the number in the MRA to the number in location n and leave the sum in the MRA.
su n (76)	Subtract from the number in the MRA the number in location n and leave the difference in the MRA.
mr n (49)	Multiply the number in the MRA by the number in location n and leave the product in the MRA.
dv n (74)	Divide the number in the MRA by the number in location n and leave the quotient in the MRA.
ts n (48)	Transfer the number in the MRA to location n.
ex n (48)	Exchange the number in the MRA with the numbers in location n.
sp m (25)	Interpret next the instruction in register m (unless m = <u>ax</u> , in which case transfer control to the register following the one which contains the <u>spax</u> so that the instruction following the <u>spax</u> is performed using the Whirlwind code).
cp m (24)	If the contents of the MRA is a negative number, proceed as in <u>sp n</u> above; if positive, ignore this instruction.
ta m (22)	Transfer the address p + 1 into the right 11 digit positions of register m, leaving the left 5 digit positions unchanged; p being the address of the most recently interpreted <u>sp</u> or effective <u>cp</u> operation.

ENTR: Operations on Real (24,6,0) Floating Point
Double Register Numbers (General Subroutine)

LSR# PA 2.2 3

Instruction Code and Operation Times:

ts 48	cp 21(+),27(-)	cs 36	cm 37
ta 22	sp 25	ad 72	mr 49
ex 48	ca 38	su 76	dv 74

Preset Parameters (to be typed in program title)

vx2/pk: k-separation in storage of two registers of number
vx/pN: N-address in storage of initial register of this
subroutine

Enter-->00	ta 179r	Set address of 1 st instruction to be interpreted	196r->25	ex 198r	"dv"
01	sp 179r		26	ts 97r	
02	(p0)	x ₁ } Multiple	27	cs 102r	} Form exponent of 2 ⁻² /x ₂
03	(p0)	x ₁ ' } register	28	ad 54r	
04	(p0)	y ₁ } accumulator	29	ts 102r	
196r-->05	sr*30	"ca"	30	cs 97r	} Form and store
06	ca ax		31	mh 97r	
13r,196r->07	ca 191r	"cs"	32	ex 198r	
08	sp 95r		33	sr *2	} $\frac{2^{-2}x_2^1}{x_2^2}$
196r->09	sp 129r	"ad"	34	dv 198r	
10	p29		35	sl *15	
196r->11	ts 97r	"su"	36	ts 151r	} Form and store $\frac{2^{-2}}{x_2^2}$
12	sp 126r		37	ca 72r	
196r->13	sp 7r	"cm"	38	dv 97r	
(170r)14	(p0)	Temporary digits storage	39	sl *15	} Form $\left(\frac{2^{-2}}{x^2}\right)^2$
24r->15	sa 3r	Add two minor products	40	ts 198r	
16	ts 3r		41	mh 97r	
17	ca 0	Store	42	su 72r	} (Use Euclid's algorithm)
18	ex 198r	overflow	43	sl *15	
19	mh 2r	Form major product	44	su 17r	
20	sp 158r		45	ad 50r	} Add two minor parts of reciprocal,
49r,196r->21	mr 2r	"mr" Form two minor products	46	dv 97r	
22	ex 3r		47	sl *15	
23	mr 198r		48	ad 151r	}
24	sp 15r		49	sp 21r	

50	pl		80r= 83	sl 14	} Add overflow into x_1 and x_2
196r= 81	<u>sp 73r</u>	"ts"	84	ts 2r	
111r= 82	ca 2r	} Complement x_1	85	ao 4r	} Increase y_1 .
53	<u>sp 164r</u>				
54	p2		82r= 86	cm 4r	} $ y_1 - 63 > 0?$
196r= 55	ca 201r	} Store digits in indicated address	87	su 62r	
(181r)56	td(0)		88	<u>cp_93r</u>	} $y_1 \leq 0?$ (i.e. $y_1 < -63?$)
57	<u>sp 178r</u>		89	cs 4r	
119r= 58	ao 2r	} Increase x_1 by 2^{-15}	90	<u>cp_202r</u>	} Set $y_1 = -63$
59	<u>sp 167r</u>		91	cs 62r	
60	sp 35r		92	ts 4r	} $ts\ n + k$ or $ex\ n + k$
196r= 61	<u>sp 73r</u>	"ex"	88r= 93	ca 97r	
62	p63		94	ad 197r	} Store $ts, ex, ca, cs,$ or $cm\ n+k$
196r= 63	cs 2r	"cp" Is x_1 negative?	8r= 95	ts 102r	
64	<u>cp_178r</u>	} Set return address for <u>sp, ax</u> .	96	ca 2r	} Perform $ts, ex, ca, cs,$ or cm
196r= 65	ao 179r		"sp"	(180r)97	
66	td 201r	} Set pick up order for ordinary cp & sp	98	ex 3r	} $x_1 \neq 0?$
67	ca 180r		99	sr *9	
68	td 179r	} Test to see whether instruction is <u>sp ax</u>	100	ex 4r	} $x_1' \neq 0?$
69	su 6r		101	sl *9	
70	<u>cp_199r</u>		(95r)102	(p0)	} Form
71	<u>sp 179r</u>		103	sr *9	
72	0.20000		104	ex 3r	
51r, 61r= 73	ca 3r	} Round off x_1' and store $x_1' \times 2^{-6}$	105	ts 2r	} $x_1' \neq 0?$
74	sr 6		106	sl *15	
75	ts 3r		107	ex 3r	} $x_1' \neq 0?$
76	sr *9	} Add round-off carry into x_1	108	<u>sp 177r</u>	
77	sa 2r		166r= 109	cm 2r	} $x_1' \neq 0?$
78	ts 2r		110	su 0	
79	ca 0	} Is there an overflow?	111	<u>cp 52r</u>	} $x_1' \neq 0?$
80	<u>cp_83r</u>		112	cm 3r	
81	su 0		113	su 0	} Form
82	<u>cp 86r</u>		114	<u>cp_122r</u>	
			115	su 50r	

OPERATIONS on Real (24, 6, 0) Floating Point
Double Register Numbers (General Subroutine)

MSR # PA 2.2 t

116	ad 17r	} $ x_1' - 1$	147	ca 50r	} Set $y_2 = 1$
117	ts 3r		148	ex 102r	
118	ca 2r	} $x_1 > 0?$	149	sr *15	} Form and store
119	cp_58r		150	ad 198r	
120	su 50r	} Form	(146r) 151	(p0)	}
121	ts 2r		152	ts 198r	
114r → 122	cs 3r	} Complement x_1'	153	sl *15	}
123	ts 3r		154	ex 3r	
124	mr 2r	} Form $x_1 \cdot x_1'$	155	sr *15	} Form
125	sp_166r		156	ad 2r	
12r → 126	cs 198r	} Complement	157	sr *1	}
127	ts 198r		20r → 158	ts 2r	
128	cs 97r	} x_2, x_2'	159	sl *15	}
9r → 129	ex 102r		160	sa 3r	
130	su 4r	} Form and store	161	ts 3r	} Add x_1, x_1'
131	ts 97r		162	ca 198r	
132	cp_141r	} $y_2 - y_1 > 0?$	163	ad 2r	}
133	ad 4r		53r → 164	ts 2r	
134	ts 4r	} Interchange	165	mr 3r	} Does sign $x_1 = \text{sign } x_1'$?
135	ca 2r		125r → 166	cp_109r	
136	ex 198r	} (x_1, x_1', y_1)	59r → 167	ca 3r	} Scale factor and
137	ts 2r		168	sr *15	
138	ca 3r	} and (x_2, x_2', y_2)	169	ad 2r	}
139	ex 102r		170	sf 14r	
140	ts 3r	}	171	ts 2r	}
132r → 141	cm 97r		172	sl *15	
142	su 10r	} $ y_2 - y_1 - 29 > 0?$	173	ts 3r	}
143	cp_145r		174	cs 14r	
144	sp_178r	} No need for addition	175	ad 102r	} Form exponent
143r → 145	ad 5r		176	ad 4r	
146	ts 151r	} Store	108r → 177	ts 4r	}
			sr *1 + $ y_2 - y_1 $		

178	no 179r	Increase address	(181r)191	ca(0)	y_2 in reg. 192
171r, 200r	ca(0)	Pick up next instruction	192	sr *9	Hold x_2' in AG.
180	ts 97r	Store instruction and digits	193	ex 102r	
181	td 56r		194	ts 198r	
182	td 189r		195	sl *15	
183	ad 197r		(188r)196	(p0)	Go to part of I.S. for particular instruction
184	td 191r		197	pe x2	
185	sr *27	Form sp to address for particular instruction	198	(p0)	Temporary storage
186	sl *17		70r= \Rightarrow 199	ad 50r	Does address equal ax?
187	ad 60r		200	sp_179	
188	ts 196r		(66r)201	(sp(0))	Return to register following <u>sp ax</u> .
(182*)189	ca(0)	Pick up x_2, x_2' and y_2 .	90r= \Rightarrow 202	ca 108r	Produce overflow
190	ts 102r	Store x_2 in reg. 198	203	ad 108r	alarm

Instructions in integrative routine: start at 200

600 | sp 179 ad
 ca 201 ad
 su 102 (ad)
 sp + data
 sp 201 ad

601 | ca
 2 | qp

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on Real (30,0,0) Fixed-Point Double Register Numbers (Minimal Routine with Sign Agreement, giving 28 Binary Digit Accuracy in mr)			LSR PA 3.1 t
			Classification Interpretive
No. of Regs. in Subroutine 96	Temp. Regs. used by Sub. -----	Average Time (operations) see description	Max. Time (operations) see description
Preset Parameters (to be inserted in main program) v8/ pk: k = separation in storage between the two registers of a double register number vx/ pN: N = address in storage of the first register of the interpretive subroutine			
Description By means of this subroutine various logical and arithmetic operations can be performed upon real (30,0,0) double-register numbers. The numbers are stored in two registers whose addresses in storage differ by a preset parameter k, i.e. if the major half of a double register number is stored in register <u>n</u> , then the minor half is stored in register <u>n+k</u> . The operations are written in the usual WW instruction code, but the meanings of these operations may differ from the usual ones (see description of instruction code). Any number of these instructions may be performed in sequence by placing an <u>sp ax</u> before the first instruction in the sequence. The instructions in the sequence are then interpreted successively until a change-of-control instruction is reached at which point either another sequence of instructions is interpreted, or, if the change-of-control instruction is an <u>sp ax</u> ordinary WW operation is resumed at the register following the instruction <u>sp ax</u> . The multiple register accumulator (mra), in which the results of instructions are left, consists of the storage registers 2r and 3r.			

TITLE: Operations on Real (30,0,0) Fixed-Point Double Register Numbers (Minimal Routine with Sign Agreement, giving 28 Binary Digit Accuracy in mra)		LSR	PA	3.1 t
<u>Instruction</u>	<u>Function</u>	<u>Av. No. Operations</u>		
ts n	Transfer the contents of mra to registers <u>n</u> and <u>n+k</u>			22
ta n	If m is the address of the last sp instruction executed by the subroutine, transfer (m+1) to the last 11 digits of register n			18
ex n	Exchange the contents of mra with the contents of <u>n</u> and <u>n+k</u>			22
cp n	If the contents of mra is negative, proceed as in the sp instruction, if the contents are positive disregard the instruction			22
sp n	If <u>sp n</u> = <u>sp ax</u> , take the next instruction to be interpreted from register n, if <u>sp n</u> = <u>sp ax</u> , resume ordinary WW operation at the register following the instruction <u>sp ax</u>			22
ca n	Clear mra, put the contents of registers <u>n</u> and <u>n+k</u> in mra			23
cs n	Clear mra, put the complement of the contents of registers <u>n</u> and <u>n+k</u> in mra			22
ad n	Add the contents of mra to the contents of registers <u>n</u> and <u>n+k</u> and store the result in mra.			40
su n	Subtract the contents of registers <u>n</u> and <u>n+k</u> from the contents of mra, store the result in mra.			44
cm n	Clear mra, put the absolute value of the contents of registers <u>n</u> and <u>n+k</u> in mra			22
mr n	Multiply the contents of mra by the contents of registers <u>n</u> and <u>n+k</u> , store the result in mra.			34

TITLE: Operations on Real (30,0,0) Fixed-Point
Double Register Numbers (Minimal Routine
with Sign Agreement, giving 28 Binary
Digit Accuracy in mr)

LSR PA 3.1 t

Notes:

1. Entering and Leaving Subroutine

Both entering and leaving the subroutine are accomplished by the instruction sp ax, where vx, a preset parameter, is the address in storage of the first register of the interpretive subroutine. When used to enter the subroutine, the first instruction interpreted is that following the instruction sp ax in storage. When used to leave the subroutine, ordinary WW operation is resumed at the register following the instruction sp ax.

2. Accuracy

All of the operations executed by the subroutine are carried out with a 30 binary digit accuracy except mr, which is carried out with 28 binary digit accuracy.

3. Sign Agreement

A sign agreement routine has been incorporated into the interpretive subroutine. This means that the major and minor halves of the number contained in the mra always have the same algebraic sign after an instruction has been executed by the interpretive subroutine.

4. Reasons for Machine Stoppage During the Subroutine

(a) Arithmetic overflow at $55r - \text{sp } n+k - 1$, i.e. an excessive address is being used for storing the minor half of a double register constant.

(b) Arithmetic overflow at 71r or 72r - overflow in sum or difference for the interpreted instructions ad or su.

(c) Various alarms can result from trying to interpret an instruction which was not meant to be interpreted by the subroutine.

5. In any of the operations the address n may be the address of the mra.

6. The mra consists of the registers 2r and 3r.

7. If $k = 1$, the two parts of the double register number are stored in consecutive storage registers.

8. No subroutine using preset parameter v8 should be used unless it is later reset.

TITLE: Operations on Real (30,0,0) Fixed-Point Double
 Register Numbers (Minimal Routine with Sign
 Agreement, giving 28 Binary Digit Accuracy in mr)

ISR# PA 3.1 t

Abstract: This subroutine is a (30,0,0) interpretive subroutine which performs the instructions ta, ta, ex, cp, sp, ca, cs, ad, su, cm and mr. The double register constants dealt with by the subroutine are in registers whose addresses in storage differ by a preset parameter k, i.e. the $C(n, n+k)$ represents a double register number. Exit and entry to the subroutine are accomplished by the instruction sp ax. In leaving the subroutine, if $C(m) = sp ax$, then ordinary WW operation is resumed at register $m+1$. In the description given below, the subroutine is assumed to be executing the instruction $C(m) = xx n$. There are two preset parameters.

Preset Parameters: (to be inserted in main program)

- v8 pk: k = the amount by which the halves of a double register number are separated in storage
- vx pN: N = the address in storage of the first register of the interpretive subroutine

00	ta	51r	} Enter interpretive subroutine	18	ex	2r	} mr
01	<u>sp</u>	51r		19	mh	28r	
02	(p0)	} mra		20	ts	28r	
03	(p0)			21	sl	15	
64r-->04	sr	19r	ca	22	<u>sp</u>	69r	
64r-->05	<u>sp</u>	44r	cs	87r--->23	ao	2r	Form $C(2r) + 2^{-15}$
64r-->06	<u>sp</u>	69r	ad	24	<u>sp</u>	50r	
64r-->07	<u>sp</u>	65r	su	79r--->25	cs	2r	Complement $C(2r)$
64r-->08	<u>sp</u>	44r	cm	26	<u>sp</u>	49r	
09	p1			64r--->27	<u>sp</u>	44r	ts
10	p0a8		v1	(62r)28	(p0)		
11	td	ax	vx	64r--->29	ca	43r	ta
64r-->12	mr	2r	}	(53r)30	td	(0)	Transfer $m+1$ to digit
13	ex	3r		31	<u>sp</u>	50r	section of register n
14	mr	28r		64r--->32	<u>sp</u>	44r	ex
15	sa	3r		64r--->33	<u>sp</u>	93r	cp
16	ts	3r		95r, 64r--->34	ao	51r	sp
17	ca	0		35	td	43r	Store $sp(m+1)$ in 43r

MINI: Operations on Real (30,0,0) Fixed-Point Double
 Register Numbers (Minimal Routine with Sign
 Agreement, giving 28 Binary Digit Accuracy in mr)

ISR# PA 3.1 t

36	ca 30r	} Store ca n in 51r	66	cs 28r	} su
37	td 51r		67	ts 28r	
38	su 11r	68	cs 45r		
39	<u>cp</u> 41r	} Is xx n = sp ax?	69	sa 3r	} ad, su, mr
40	<u>sp</u> 51r		70	ts 3r	
39r-41	ad 9r	71	ca 2r		
42	<u>cp</u> 51r	72	ad 28r		
(35r)43	<u>sp</u> (0)	73	ts 2r		
5r, 8r, } 27r, 32r } (56r)	44 ca 3r	} Perform	92r-94	cs 2r	} Do C(2r) and C(3r) disagree in sign?
45	(p0)		75	mr 3r	
46	ts 3r	} ts, ex, ca, cs, cm	76	<u>cp</u> 50r	} Is C(2r) ≠ 0?
47	ca 2r		77	cm 2r	
(54r)48	(p0)	78	su 0		
26r-49	ts 2r	} Is C(3r) ≠ 0?	79	<u>cp</u> 25r	
24r, 31r } 76r, 94r } (Cr), 1r } 40r, 42r }	50 ao 51r		80	cm 3r	
51	ca (0)	81	su 0		
52	td 61r	} Store n	82	<u>cp</u> 90r	} Form 1 - C(3r)
53	td 30r		83	su 9r	
54	ts 48r	84	ad 17r		
55	ad 10r	85	ts 3r		
56	ts 45r	86	ca 2r	} Is C(2r) positive?	
57	td 63r	87	<u>cp</u> 23r		
58	sr *11	} Set up entry into table	88	su 9r	} Form C(2r) - 2 ⁻¹⁵
59	ad 4r		89	ts 2r	
60	td 64r	82r-90	cs 3r	} Complement C(3r)	
(52r)61	ca (0)	91	ts 3r		
62	ts 28r	92	<u>sp</u> 74r	Re-enter sign agreement	
(57r)63	ca (0)	33r-93	cs 2r	} cp	
(60r)64	<u>sp</u> (0)	94	<u>cp</u> 50r		} Is C(2r) negative?
7r-65	ts 45r	95	<u>sp</u> 34r	Perform sp order.	

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on Real (30,0,0) Fixed-Point Double-Register Numbers (General Routine with Sign Agreement, No Division) (Interpretive)			LSR PA 3.5t
			TAPE T721-2
			Classification Interpretive
No. of Regs. in Subroutine 123	Temp. Regs. used by Subroutine -----	Average Time (operations) see description	Max. Time (operations) see description
Preset Parameters vx2/ pk: k = separation in storage between the two registers of a double register number vx/ pN: N = address in storage of the first register of the interpretive subroutine			
Description By means of this subroutine various logical and arithmetic operations can be performed upon real (30,0,0) double-register numbers. The numbers are stored in two registers whose addresses in storage differ by a preset parameter k, i.e. if the major half of a double register number is stored in register <u>n</u> , then the minor half is stored in register <u>n+k</u> . The operations are written in the usual WW instruction code, but the meanings of these operations may differ from the usual ones (see description of instruction code). Any number of these instructions may be performed in sequence by placing an <u>sp ax</u> before the first instruction in the sequence. The instructions in the sequence are then interpreted successively until a change-of-control instruction is reached at which point either another sequence of instructions is interpreted, or, if the change-of-control instruction is an <u>sp ax</u> ordinary WW operation is resumed at the register following the instruction <u>sp ax</u> . The multiple register accumulator (mra), in which the results of instructions are left, consists of the storage registers 2r and 3r.			
<u>Instruction</u>	<u>Function</u>	<u>Av. No. Operations</u>	
ts n	Transfer the contents of mra to registers <u>n</u> and <u>n+k</u>	22	
ta n	If m is the address of the last sp instruction executed by the subroutine, transfer (m+1) to the last 11 digits of register n	18	
ex n	Exchange the contents of mra with the contents of <u>n</u> and <u>n+k</u>	22	
cp n	If the contents of mra is negative, proceed as in the sp instruction, if the contents are positive disregard the instruction	22	

TITLE: Operations on Real (30,0,0) Fixed-Point Double-Register Numbers (General Routine with Sign Agreement, No Division) (Interpretive)

LSR PA 3.5t

<u>Instruction</u>	<u>Function</u>	<u>Av. No. Operations</u>
sp n	If $sp\ n = sp\ ax$, take the next instruction to be interpreted from register n , if $sp\ n \neq sp\ ax$, resume ordinary WW operation at the register following the instruction $sp\ ax$	22
ca n	Clear mra, put the contents of registers n and $n+k$ in mra	23
es n	Clear mra, put the complement of the contents of registers n and $n+k$ in mra	22
ad n	Add the contents of mra to the contents of registers n and $n+k$ and store the result in mra	40
su n	Subtract the contents of registers n and $n+k$ from the contents of mra, store the result in mra	44
cm n	Clear mra, put the absolute value of the contents of registers n and $n+k$ in mra	22
mr n	Multiply the contents of mra by the contents of registers n and $n+k$, store the result in mra.	34
sl(800+n)	Multiply C(mra) by 2^n and leave the result in the mra	29
sr(800+n)	Multiply C(mra) by 2^{-n} , and store the first 30 digits of the result without roundoff in the mra.	28

Notes

1. Entering and Leaving Subroutine

Both entering and leaving the subroutine are accomplished by the instruction $sp\ ax$, where vx , a preset parameter, is the address in storage of the first register of the interpretive subroutine. When used to enter the subroutine, the first instruction interpreted is that following the instruction $sp\ ax$ in storage. When used to leave the subroutine, ordinary WW operation is resumed at the register following the instruction $sp\ ax$.

2. Accuracy

All of the operations executed by the subroutine are carried out with 30 binary digit accuracy, affected only by roundoff on the thirty-first digit.

3. Sign Agreement

A sign agreement routine has been incorporated into the interpretive subroutine. This means that the major and minor halves of the number contained in the mra always have the same algebraic sign after an instruction has been executed by the interpretive subroutine.

4. Reasons for machine stoppage during the subroutine

(a) Arithmetic overflow at $56r = sp\ n+k$, an excessive address is being used for storing the minor half of a double register constant.

TITLE: Operations on Real (30,0,0) Fixed-Point Double-
Register Numbers (General Routine with Sign
Agreement, No Division) (Interpretive)

LSR PA 3.5t

(b) Arithmetic overflow at 72r or 73r - overflow in sum or difference for the interpreted instructions ad or su.

(c) Various alarms can result from trying to interpret an instruction which was not meant to be interpreted by the subroutine.

5. The mra consists of registers 2ax (at 2r) and 3ax (at 3r).

FCH 12/5/51	AS 2/5/52
----------------	--------------

TITLE: Operations on Real (30,0,0) Fixed-Point
 Double-Register Numbers (General Routine
 with Sign Agreement, No Division)
 (Interpretive)

LSR# PA 3.5t
 TAPE# T721-2

Abstract: This subroutine is a (30,0,0) interpretive subroutine which performs the instructions ts, ta, ex, cp, ca, cs, ad, su, cm, mr, sl, and sr. The double register constants dealt with by the subroutine are in registers whose addresses in storage differ by a preset parameter k, i.e. the $C(n, n+k)$ represents a double register number. Exit and entry to the subroutine are accomplished by the instruction sp ax. In leaving the subroutine, if $C(m) = sp ax$, then ordinary WW operation is resumed at register $m+1$. In the description given below, the subroutine is assumed to be executing the instruction $C(m) = xx n$. There are two preset parameters.

Preset Parameters: (to be inserted in main program)

vx2: pk: k = the amount by which the halves of a double register number are separated in storage
 vx: pN: N = the address in storage of the first register of the interpretive subroutine

00	ta 51r	} Set next instruction address	20	ts 2r	}
01	<u>sp 51r</u>		21	sl* 15	
(23r) → 02	(p0)	} mra	22	<u>sp 92r</u>	}
03	(p0)		88r → 23	ao 2r	
65r → 04	sr 19r	ca	24	<u>sp 50r</u>	
65r → 05	<u>sp 44r</u>	os	80r → 25	cs 2r	Complement $C(2r)$
65r → 06	<u>sp 70r</u>	ad	26	<u>sp 49r</u>	
65r → 07	<u>sp 66r</u>	su	65r → 27	<u>sp 44r</u>	ts
65r → 08	<u>sp 44r</u>	cm	(68r)(63r) 28	(p0)	$C(n) = x_2$
33r → 09	cs 2r	} Is $C(2r)$ negative?	65r → 29	ca 43r	ta
10	<u>op 50r</u>		(53r) 30	td(0)	} Transfer $m+1$ to digit section of register D
11	<u>sp 34r</u>	Perform sp instruction	31	<u>sp 50r</u>	
65r → 12	ts 19r	} mr	65r → 32	<u>sp 44r</u>	ex
13	mr 3r		65r → 33	<u>sp 9r</u>	op
14	<u>sp 94r</u>		11r, 65r → 34	ao 51r	sp
65r → 15	ca ax	sl (vx)	35	td 43r	} store $sp(m+1)$ in 43r
65r → 16	ca 3r	sr	36	ca 62r	
17	sr* 15	} sl, sr $C(n+k)=x_2^1$	37	td 51r	} Is $xxn = spax$
18	ad 2r		38	su 15r	
(103r)(55r)(12r) → 19	(p0)		39	<u>cp 41r</u>	

TITLE: Operations on Real (30,0,0) Fixed-Point
 Double-Register Numbers (General Routine
 with Sign Agreement, No Division) (Interpretive)

LSR# PA 3.5t
 TAPE# T721-2

	40	<u>sp 51r</u>		6r	→	70	sa 3r	
39r →	41	ad 122r		120r		71	ts 3r	
	42	<u>op 51r</u>				72	ca 2r	} ad, su, mr
(35r)	43	<u>sp 0</u>	} Leave interpretive subroutine			73	ad 28r	
5r, 8r	44	ca 3r					74	
27r, 32r	45	(p0)	} Perform ts, ex, ca, cs, and cm.	93r →	75	cs 2r	} Do C(2r) and C(3r) agree in sign	
(97r)(66r)(57r)	46	ts 3r				76		mr 3r
(113r)	47	ca 2r				77	<u>op 50r</u>	
	48	(p0)				78	cm 2r	} Is C(2r) ≠ 0?
(54r)	49	ts 2r			79	su0		
26r →	50	ac 51r	} Pick up next instruc- tion to be interpreted		80	<u>cp 25r</u>	} Is C(3r) ≠ 0?	
10r, 24r	51	ca(0)				81		cm 3r
31r, 77r	52	td 62r	} Store n		82	su0	} Is C(3r) ≠ 0?	
1r, 40r, 42r, (0r)	53	td 30r				83		<u>op 91r</u>
(34r)(37r)(50r)	54	ts 48r	} Store xxn		84	su 122r	} Form -1 + C(3r)	
	55	ts 19r				85		ad 118r
	56	ad 121r			86	ts 3r	} Is C(2r) positive?	
	57	ts 45r	Store xx(n+k)		87	ca 2r		
	58	td 64r	Store (n+k)		88	<u>op 23r</u>	} Form C(2r) - 2 ⁻¹⁵	
	59	sr* 11	} Set up entry into table		89	su 122r		
	60	ad 4r				90	ts 2r	
	61	td 65r			83r →	91	cs 3r	} Complement C(3r)
(52r)	62	ca(0)	} Store C(n) in 28r	22r →	92	ts 3r		
	63	ts 28r				93	<u>sp 75r</u>	Re-enter sign agreement
(58r)	64	ca(0)	put C(n+k) = x ₂ ' in	AC 14r →	94	sr 2	} Form $\frac{1}{4} x_1' x_2'$ and store in 19r	
(61r)	65	<u>sp(0)</u>	enter table		95	ex 19r		
7r →	66	ts 45r	} su		96	mh 2r	} Form x ₁ x ₂ ', store major half in 45r	
	67	cs 28r				97		ts 45r
	68	ts 28r				98	sl*15	} Form $\frac{1}{4} [x_1' x_2']$
	69	cs 45r				99	sr 2	

TITLE: Operations on Real (30,0,0) Fixed-Point
 Double-Register Numbers (General Routine
 with Sign Agreement, No Division)(Interpretive)

LSR# PA 3.5t
 TAPE # T721-2

100	ad 19r	} Accumulate round-off and store in 2r	112	sa 45r	} Add $x_1 x_2'$ and the major half of $x_1 x_2'$
101	ex 2r		113	ts 45r	
102	mh 28r	} Form $x_1 x_2$, store major half in 19r and minor half in 28r	114	ca 19r	
103	ts 19r		115	ex 28r	
104	sl* 15		116	sa 45r	
105	ex 28r	} Form $x_1' x_2$ and store major half in 3r	117	ts 2r	
106	mh 3r		118	ca0	
107	ts 3r	} Form $\frac{1}{4} [x_1' x_2']$	119	ex 2r	
108	sl* 15		120	<u>sp 70r</u> Enter addition routine	
109	sr 2	} Form total round-off and add into major half of $x_1 x_2'$	121	p0ax2 v1	
110	ad 2r		122	p1	
111	sr 13				

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on Real (30,0,0) Fixed-Point Double Register Numbers (Short, fast routine without sign agreement giving 28 binary digit accuracy in mr) Interpretive		LSR PA 3.10t <hr/> TAPE T798 <hr/> Classification Interpretive	
No. of Regs. in Subroutine 78	Temp. Regs. used by Sub. 0	Average Time (operations) -----	Max. Time (operations) -----
Preset Parameters (To be inserted in title of main program) vx pN, where N is the address in storage of the first register of the interpretive subroutine vx2 pk, where k is the separation in storage between the two registers of a double register number			
Description By means of this subroutine various logical and arithmetic operations can be performed upon real (30,0,0) double-register numbers. The numbers are stored in two registers whose addresses in storage differ by a preset parameter k, i.e. if the major half of a double register number is stored in register <u>n</u> , then the minor half is stored in register <u>n+k</u> . The operations are written in the usual WW instruction code, but the meanings of these operations may differ from the usual ones (see description of instruction code). Any number of these instructions may be performed in sequence by placing an <u>sp ax</u> before the first instruction in the sequence. The instructions in the sequence are then interpreted successively until a change-of-control instruction is reached at which point either another sequence of instructions is interpreted, or, if the change-of-control instruction is an <u>sp ax</u> ordinary WW operation is resumed at the register following the instruction <u>sp ax</u> . The multiple register accumulator (MRA), in which the results of instructions are left, consists of the storage registers 2r and 3r. This interpretive subroutine does not contain a sign agreement routine and hence the contents of the MRA may have different signs, e.g. in the subtraction $10 \cdot 2^{-15} + 5 \cdot 2^{-30} - [5 \cdot 2^{-15} + 10 \cdot 2^{-30}] = 5 \cdot 2^{-15} + (-10) \cdot 2^{-30}$ the result has major and minor halves of unlike sign. This fact must be remembered by the programmer if computations are carried on outside of the subroutine or if other library subroutines are interpreted by the interpretive subroutine. In the latter case, a library subroutine should not be used unless it is explicitly stated in its specification sheet that it is permissible to use it with interpretive subroutines not containing sign agreement.			

TITLE: Operations on Real (30,0,0) Fixed-Point Double
 Register Numbers (Short, fast routine without
 sign agreement giving 28 binary digit accuracy
 in mr) Interpretive

LSR PA 3.10t
 TAPE T798

<u>Instruction</u>	<u>Function</u>	<u>Av. No. Operations</u>
ts	Transfer the contents of <u>MRA</u> to registers <u>n</u> and <u>n+k</u>	17
ta	If <u>m</u> is the address of the last sp instruction executed by the subroutine, transfer (<u>m+1</u>) to the last 11 digits of register <u>n</u>	13
ex	Exchange the contents of <u>MRA</u> with the contents of <u>n</u> and <u>n+k</u>	17
cp	If the contents of <u>MRA</u> is negative, proceed as in the sp instruction, if the contents are positive disregard the instruction	(+)14 (-)25
sp	If <u>sp n</u> \neq <u>sp ax</u> , take the next instruction to be interpreted from register <u>n</u> , if <u>sp n</u> = <u>sp ax</u> , resume ordinary WW operation at the register following the instruction sp ax	18
ca	Clear <u>MRA</u> , put the contents of registers <u>n</u> and <u>n+k</u> in <u>MRA</u>	17
cs	Clear <u>MRA</u> , put the complement of the contents of registers <u>n</u> and <u>n+k</u> in <u>MRA</u>	17
ad	Add the contents of <u>MRA</u> to the contents of registers <u>n</u> and <u>n+k</u>	18
su	Subtract the contents of registers <u>n</u> and <u>n+k</u> from the contents of <u>MRA</u>	17
mr	Multiply the contents of <u>MRA</u> by the contents of registers <u>n</u> and <u>n+k</u> , store the result in <u>MRA</u>	33

TITLE: Operations on Real (30,0,0) Fixed-Point Double
Register Numbers (Short, fast routine without
sign agreement giving 28 binary digit accuracy
in mr) Interpretive

LSR PA 3.10t
TAPE T798

Notes:

1. Entering and Leaving Subroutine

Both entering and leaving the subroutine are accomplished by the instruction sp ax, where vx, a preset parameter, is the address in storage of the first register of the interpretive subroutine. When used to enter the subroutine, the first instruction interpreted is that following the instruction sp ax in storage. When used to leave the subroutine, ordinary WW operation is resumed at the register following the instruction sp ax.

2. Accuracy

All of the operations executed by the subroutine are carried out with a 30 binary digit accuracy except mr, which is carried out with 28 binary digit accuracy.

3. Sign Agreement

PA 3.10t does not contain sign agreement. Hence the major and minor halves of a number need not have the same sign.

4. Reasons for Machine Stoppage During the Subroutine

(a) Arithmetic overflow at 56r--an excessive address is being used for storing the minor half of a double register number.

(b) Arithmetic overflow at 49r or 50r--overflow in sum or difference for the interpreted instructions ad or su.

(c) Various alarms can result from trying to interpret an instruction which was not meant to be interpreted by the subroutine.

5. The MRA consists of registers 2ax(2r) and 3ax(3r).

FCH
2/5/52

MAS
2/8/52

Abstract: This subroutine is a (30,0,0) interpretive subroutine which performs the instructions ts, ta, ex, cp, sp, ca, cs, ad, su, cm and mr. The double register constants dealt with by the subroutine are in registers whose addresses in storage differ by a preset parameter k, i.e. the C(n, n+k) represents a double register number. Exit-and entry to the subroutine are accomplished by the instruction sp ax. In leaving the subroutine, if C(m) = sp ax, then ordinary NW operation is resumed at register m+1. In the description given below, the subroutine is assumed to be executing the instruction C(m) = xx n. There are two preset parameters. The subroutine does not contain sign agreement. Hence the major and minor halves of a double register number can have unlike signs.

Preset Parameters

- vx pN, where N is the address in storage of the first register of the interpretive subroutine
- vx2 pK, where K is the separation in storage between the two registers of a double register number

00	ta 53r	} Enter interpretive subroutine	20	su 0	}	
01	<u>sp 53r</u>		21	<u>cp 60r</u>		
(76r)02	(p0)	} mra	22	cs 2r	} Is C(2r) or C(3r) pos.?	
(74r)03	(p0)		64r → 23	<u>cp 52r</u>		
59r → 04	<u>sp 44r</u>	ca	24	<u>sp 34r</u>	Perform interpreted sp	
59r → 05	<u>sp 44r</u>	cs	25	pa x2		
59r → 06	<u>sp 7r</u>	ad	26	1.67777ax2		
6r, 59r → 07	ad 26r	} su	50r → 27	<u>sp 44r</u>	ts	
08	ts 9r		Form ca(n+k) or cs (n+k) & store in 9r	28	sp ax	
77r → (8r)(15r)09	(p0)	} {Form $x_2^1 + x_1^1$ and store in 3r(ad, su)	59r → 29	td 65r	ta	} Transfer n to digit section of reg. 65r
10	sa 3r		30	ca 42r		
11	<u>sp 48r</u>	} {Form $x_1^1 x_2^1 + x_1^1 x_2^1$ + $x_0^1 x_2^1$ (mr)	31	<u>sp 65r</u>		
39r → 12	td 68r		mr	59r → 32	<u>sp 44r</u>	ex
13	ts 72r	} Set address at 68r to n & store mrn in 72r	59r → 62r → 33	<u>sp 19r</u>	cp	
14	ad 25r		24r → 34	ao 53r	sp	
15	ts 9r	} Store mr(n+k) in 9r	59r	35	td 42r	Store sp (m+1) in 42r
16	ca 51r		} Store ts2r in 50r	36	ca 50r	} Store ca n in 53r
17	ts 50r	37		td 53r		
18	<u>sp 67r</u>		38	su 28r	} Is xxn = sp ax?	
33r → 19	cm 2r	} Is C(2r) ≠ 0	39	ts 47r		

TITLE: Operations on Real (30,0,0) Fixed-Point Double
 Register Numbers (Short, fast routine without
 sign agreement giving 28 binary digit accuracy
 in mr) Interpretive

LSR PA 3.10t
 TAPE T798

40 cm 47r
 41 su 0
 (35r) 42 cp(0)
 43 sp 53r
 4r,5r → 44 ad 25r } Store xx(n+k) in 47r
 27r,32r → 45 ts 47r }
 46 ca 3r
 (39r) (69r) 47 (p0)
 (45r)
 11r → 48 ts 3r } Perform ts,ex,ca,cs,
 49 ca 2r } ad,su,mr
 (17r)(54r) 50 (p0)
 51 ts 2r
 23r,66r → 52 ao 53r } Pick up instruction
 (34r)(37r)(0r) 53 ca(0) } to be interpreted
 1r,43r → 54 ts 50r } Store xxn
 (52r)
 55 sr* 11 }
 56 ad 33r } Set up entry into
 57 ts 59r } table
 58 ca 50r } Pick up xxn

(57r) 59 (p0) Enter table
 21r → 60 cm 3r }
 61 su 0 } Is C(3r) ≠ 0
 62 cp 34r
 63 cs 3r
 64 sp 23r
 (29r) 65 td (0) } Transfer m+1 to digit section
 31r → 66 sp 52r } of register n
 18r → 67 ca 2r } Form x_1x_2
 68 mh(0)n
 69 ts 47r } Store $[x_1x_2]$ in 47r
 70 sl 15 } Store $[x_1x_2]^1$ in 3r
 71 ex 3r }
 (13r) 72 (p0) } Form $[x_1x_2]^1$
 73 sa 3r } Form $[x_1x_2]^1 + [x_1x_2]^1$
 74 ts 3r } and store in 3r
 75 ca 47r } Store $[x_1x_2]$ in 47r
 76 ex 2r }
 77 sp 9r

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on real (15,0,c) fixed-point single register numbers (Interpretive)			LSR PA 8.1t
			TAPE T724-2
			Classification Interpretive
No. of Regs. in Subroutine 47	Temp. Regs. used by Sub. None	Average Time (operations) See description of order code.	Max. Time (operations)
Preset Parameters (To be typed in the title of the Main Program)			
vx: pN: N = address in storage of initial register of subroutine vx3: pc: c = the number of binary digits to the right of the binary point in the (15,0,c) numbers vx1: pM: M = address in storage of the initial register of the storage block for (15,0,c) numbers			
Description By means of this subroutine various logical and arithmetic operations can be performed on real numbers expressed in the (15,0,c) system, $0 \leq c \leq 15$. The (15,0,c) constants are stored in the following manner. Let c be fixed and let a be a positive number such that $2^{-(15-c)} \leq a \leq 2^{c+1} - 1$ Then a can be written as the sum of a c digit binary integer and a binary fraction. The binary fraction is then rounded off to 15 - c digits and the result stored with the sign digit zero. If a < 0, repeat this procedure for -a, complement this number and store the result. For example, let c = 3 and a = $-3\frac{1}{3}$			
Then $-a = 3 + \frac{1}{3}$ and $3 = +.011$ $\frac{1}{3} = +.010101010101010$ <p style="text-align: center; margin-left: 100px;"><small>15-3=12 digits</small></p> $3 + \frac{1}{3} = 011.0101010101$ We now complement the following number 0.011010101010101 and store the following result 1.100101010101010			
The programmer need not carry out this conversion process himself but instead need only write the fixed point decimal number. By giving this information to the proper subroutine in the IP section of the library the conversion can be done automatically.			

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on real (15,0,c) fixed-point single register numbers (Interpretive)	LSR PA 8.1t
	TAPE T724-2

Description (continued)

Operations upon numbers are written in the usual Whirlwind instruction code, but the meaning of these operations may differ from the usual ones. Any number of these operations may be performed in sequence by placing an spax before the first instruction in the sequence. The instructions in the sequence are then interpreted successively until a change-of-control instruction is reached (see description of order code) at which point either another sequence of instructions is interpreted, or, if the change of control instruction is an spax ordinary Whirlwind operation is resumed at the register following the instruction spax.

The multiple register accumulator, in which the results of interpreted instructions are left, consists of the storage register 2ax.

<u>Inst.</u>	<u>Function</u>	<u>Average # Operations</u>
ts n	Transfer C(mra) to register n.	14
td n	Transfer the last 11 digits from the mra to the last 11 digits of register n	14
ta n	If m is the address of the last sp instruction or effective cp instruction executed by the subroutine, transfer (m+1) into the last 11 digits of register n	13
ex n	Exchange C(mra) with C(n)	14
cp n	If C(mra) is negative, proceed as in the sp instruction, if C(mra) is positive, disregard the instruction	18
sp n	If spn ≠ spax, take the next instruction to be interpreted from register n, if spn = spax, resume ordinary Whirlwind operation at the register following the instruction spax	21
ca n	Clear mra, put C(n) in the mra	14
cs n	Clear mra, put the complement of C(n) in the mra	14
ad n	Add C(mra) to C(n) and store the result in the mra	14
su n	Subtract C(n) from C(mra) and store the result in the mra	14
cn n	Clear the mra, and put the positive magnitude of C(n) in the mra	14
mr n	Multiply C(mra) by C(n), store the result in the mra	14
sl n	Multiply C(mra) by 2 ⁿ and store the result in the mra	14
sr n	Multiply C(mra) by 2 ⁻ⁿ and store the result in the mra	14

DIGITAL COMPUTER LABORATORY
MASSACHUSETTS INSTITUTE OF TECHNOLOGY
WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Operations on real (15,0,c) fixed-point single register numbers (Interpretive)	LSR PA 8.1t
	TAPE T724-2

Notes:

1. Entering and leaving subroutine

Both entering and leaving the subroutine are accomplished by the instruction sp ax, where vx, a preset parameter, is the address in storage of the first register of the interpretive subroutine. When used to enter the subroutine, the first instruction interpreted is that following the instruction sp ax in storage. When used to leave the subroutine, ordinary WW operation is resumed at the register following the instruction sp ax.

2. Accuracy

All of the operations executed by the subroutine are carried out with a 15 binary digit accuracy.

3. Calculation with integers

If c = 15, the (15,0,c) numbers dealt with by the subroutine are binary integers.

4. Reasons for machine stoppage during the subroutine

- a) Arithmetic overflow at 35r - overflow in a ~~sum~~ difference for the interpreted instructions ad or su.
- b) Various alarms can result from trying to interpret an instruction which was not meant to be interpreted by the subroutine.

5. Overflow during mr

It should be noted that overflows can occur during multiplication of (15,0,c) numbers. These will not cause an alarm since the overflow is shifted out of the left hand end of the accumulator, i.e., the product is formed modulo 2^c .

6. In any of the operations the address n may be the address of the mra.

7. The mra consists of register 2r (or 2ax).

FCH 12/6/51	AS 2/5/52
----------------	--------------

TITLE: Operations on real (15,0,c) fixed-point
single register numbers (Interpretive)

LSR# PA 8.1t
TAPE# T 724-2

Abstract: PA 8.1 is a (15,0,c) interpretive subroutine which performs the instructions ts, td, ta, ex, cp, sp, ca, cs, ad, su, cm, mr, sl and sr. Exit and entry to the subroutine are accomplished by the instruction spax. In leaving the subroutine, if C(m) = spax, then ordinary WW operation is resumed at register m + 1.

Preset Parameters:

- vx1: M address in storage of the initial register of the storage block for (15,0,C) numbers.
- vx3: C where C is the number of binary digits to the left of the binary point
- vx N where N is the address in storage of the first register of the interpretive routine

In the description given below the subroutine is assumed to be executing the order C(m) = xxn

00	ta 38r	} Enter interpretive subroutine	46r → 25	<u>sp 38r</u>		
01	<u>sp 38r</u>			26	<u>sp 34r</u>	ts
02	(p0)	mra	46r → 27	<u>sp 34r</u>	td	
46r → 03	sr 18r	ca	46r → 28	ca 24r	ta	
46r → 04	ca ax	cs	(41r) 29	td(0)	} Transfer m+1 to digit section of register n	
46r → 05	<u>sp 34r</u>	ad		30 <u>sp 37r</u>		
46r → 06	<u>sp 34r</u>	su	46r → 31	<u>sp 34r</u>	ex	
46r → 07	<u>sp 34r</u>	cm	46r → 32	<u>sp 8r</u>	cp	
32r → 08	cs 2r	} cp	46r → 33	<u>sp 16r</u>	sp	
09	<u>cp 37r</u>		} 5r,6r,7r } 14r,15r } 26r,27r,31r } (21r)(39r)	→ 34	ca 2r	} Perform ca, cs, ad, su, cm, sl, sr, ts, td, and ex
10	<u>sp 16r</u>			→ 35	(p0)	
46r → 11	mh 2r	13r → 36	ts 2r			
12	sl0ax3	} mr	9r,30r → 37	ao 38r		
13	<u>sp 36r</u>		(0r)(16r) } 1r, 25r } (37r)(19r)	38	ca(0)	
46r → 14	<u>sp 34r</u>	sl	39	ts 35r	Store xxn	
46r → 15	<u>sp 34r</u>	sr		40	td 45r	} Store n
10r,33r → 16	ao 38r	} store cp(m+1) in 24r		41	td 29r	
17	td 24r				42	sr* 11
18	ca 45r	} store ca n in 38r		43	ad 3r	
19	td 38r				44	td 46r
20	su 4r	} Is xxn=spax?	(40r) 45	ca(0)	Put C(n) in AC	
21	ts 35r		(44r) 46	<u>sp(0)</u>	Enter table.	
22	cm 35r					
23	su0					
(17r) 24	<u>cp0</u>					

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Form $\sin \frac{\pi}{2} x$ from x Stored in AC, and Leave Result in AC			LSR TF 0.1 t
			Classification Closed
No. of Regs. in Subroutine	Temp. Regs. used by Sub.	Average Time (operations)	Max. Time (operations)
19	d - 1t	15	15
Program Parameters on entering Subroutine ac: x ar: return address			
Results on leaving Subroutine ac: $\sin \frac{\pi}{2} x$			
Description The subroutine gives $\sin \frac{\pi}{2} x = ax - bx^3 + cx^5 - dx^7$ for $-1 < x < 1$. The subroutine is entered with x in the accumulator and on returning to the main program, $\sin \frac{\pi}{2} x$ is in the AC. The maximum error is approximately ± 0.00005 and the average error is ± 0.00003 .			

TITLE: Form $\sin \frac{\pi}{2} x$ from x Stored in AC, and Leave

LSR# TF 0,1 t

Result in AC.

Abstract: This subroutine gives $\sin \frac{\pi}{2} x = ax - bx^3 + cx^5 - dx^7$ as a single length fixed point number in the accumulator where $-1 < x < 1$.

Entering the subroutine:

ac: x

Leaving the subroutine:

ac: $\sin \frac{\pi}{2} x$

Temporary Storage:

d unused

lt used to store the value of x

00	ta 14r	Set return address	10	sr *1	$(-bx^2 + cx^4 - dx^6)2^{-1}$
01	ts lt	Store x	11	ad 18r	$(a - bx^2 + cx^4 - dx^6)2^{-1}$
02	mh lt	x^2	12	mh lt	$(ax - bx^3 + cx^5 - dx^7)2^{-1}$
03	mh 15r	$-dx^2$	13	sl 1	$ax - bx^3 + cx^5 - dx^7$
04	ad 16r	$c - dx^2$	(Or) 14	<u>sp (0)</u>	Return to main program
05	mh lt	$cx - dx^3$	15	1,77560	-d
06	mh lt	$cx^2 - dx^4$	16	0,05055	c
07	ad 17r	$-b + cx^2 - dx^4$	17	1,26521	-b
08	mh lt	$-bx + cx^3 - dx^5$	18	0,62210	$a x 2^{-1}$
09	mh lt	$-bx^2 + cx^4 - dx^6$			

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Form $\cos \frac{\pi}{2} y$ from y Stored in AC, and Leave Result in AC (15,0,0)			LSR TF 1.1t
			Tape 705-1
			Classification Closed
No. of Regs. in subroutine 26	Temp. Regs. used by Sub. d -1t(2)	Average Time (operations) 20	Max. Time (operations) 20
Program Parameters on entering Subroutine ac: y ar: Return address			
Results on leaving Subroutine ac: $\cos \frac{\pi}{2} y$			
Description Gives $\cos \frac{\pi}{2} y$ by changing y to x so that $\cos \frac{\pi}{2} y = \sin \frac{\pi}{2} x = ax - bx^3 + cx^5 - dx^7$, for $-1 < x < 1$. If $y=0$, there will be an overflow since the cosine of zero is one. The sub- routine is entered with y in the accumulator and on returning to the main program $\cos \frac{\pi}{2} y$ is in the accumulator. The maximum error is approximately ± 0.00005 and the average error ± 0.00002 .			
Notes 1. If y is <u>zero</u> , an arithmetic overflow will occur in register 5r.			
DMN+DGA 1/14/52	JWC YII 1/16/52	MAS 1/22/52	

TITLE: Form cosine $\frac{\pi}{2}y$ from y Stored in AC, and
 Leave Result in AC (15,0,0)

LSR# TF 1.1t
 Tape 705-1

Abstract: Gives cosine $\frac{\pi}{2}y$ by forming $x=1-|y|$ so that
 cosine $\frac{\pi}{2}y = \text{sine } \frac{\pi}{2}x$ and evaluating
 $\text{sine } \frac{\pi}{2}x = ax - bx^3 + cx^5 - dx^7$, $-1 < x < 1$. If $y = \underline{+0}$
 there will be an overflow at register 5r.

Upon entering the subroutine:
 ac: y

Upon leaving the subroutine:
 ac: cosine $\frac{\pi}{2}y$

Temporary registers:
 d - unused
 lt - used to store x

00	ta 19r	Set return address	13	mh lt	$-bx + cx^3 - dx^5$
01	ts lt	Transfer y to lt	14	mh lt	$-bx^2 + ox^4 - dx^6$
02	cp 4r	y negative? Yes	15	sr*1	$(-bx^2 + ox^4 - dx^6)2^{-1}$
03	os lt	No. -y in ac	16	ad 25r	$(a-bx^2 + cx^4 - dx^6)2^{-1}$
04	ad 20r	$(1 - 2^{-15}) -y$	17	mh lt	$(ax - bx^3 + cx^5 - dx^7)2^{-1}$
05	ad 21r	Add 2^{-15} } (Possible overflow)	18	sl 1	$ax - bx^3 + cx^5 - dx^7$
06	ts lt	$1 - y = x$ in it	(Or)19	sp(0)	Return to main program
07	mh lt	x^2	20	0.77777	$1 - 2^{-15}$
08	mh 22r	$-dx^2$	21	0.00001	2^{-15}
09	ad 23r	$c - dx^2$	22	1.77560	-d
10	mh lt	$ox - dx^3$	23	0.05055	c
11	mh lt	$ox^2 - dx^4$	24	1.26521	-b
12	ad 24r	$-b + ox^2 - dx^4$	25	0.62210	$a x 2^{-1}$

DIGITAL COMPUTER LABORATORY
 MASSACHUSETTS INSTITUTE OF TECHNOLOGY
 WHIRLWIND SUBROUTINE SPECIFICATION

TITLE: Form Sine $\frac{\pi}{2}x$ from x Stored in AC, and/or Form Cosine $\frac{\pi}{2}y$ from y Stored in AC, Leave Result in AC, (15,0,0)			LSR TF 7.1 t
			Tape T 750-1
			Classification Closed
No. of Regs. in Subroutine 28	Temp. Regs. used by Sub. d - 1t	Average Time (operations) 15 - sine 20 - cosine	Max. Time (operations) 15 - sine 20 - cosine
Program Parameters on entering Subroutine ac: x or y ar: return address			
Results on leaving Subroutine ac: sine $\frac{\pi}{2}x$ or cosine $\frac{\pi}{2}y$			
Description If this subroutine is entered at register Or, it will calculate cosine $\frac{\pi}{2}y$ by changing y to x so that cosine $\frac{\pi}{2}y = \text{sine } \frac{\pi}{2}x$ and evaluating $\text{sine } \frac{\pi}{2}x = ax - bx^3 + cx^5 - dx^7, -1 < x < 1$. If it is entered at register 7r, it will calculate sine $\frac{\pi}{2}x$. There will be an overflow at register 5r if y = 0. The subroutine is entered with either x or y in the accumulator, and on returning to the main program either sine $\frac{\pi}{2}x$ or cosine $\frac{\pi}{2}y$ is in the accumulator. The maximum error is approximately +0.00005 and the average error is +0.00002.			
Notes: 1. Enter at Or if cosine $\frac{\pi}{2}y$ is desired enter at 7r if sine $\frac{\pi}{2}x$ is desired 2. If y = +0, an overflow will occur in register 5r.			
DMN/DGA	JWC III	DEL	
1/14/52	1/16/52	1/18/52	

TITLE: FORM SINE $\frac{\pi}{2}x$ FROM x STORED IN AC, AND/OR
 FORM COSINE $\frac{\pi}{2}y$ FROM y STORED IN AC,
 LEAVE RESULT IN AC, (15,0,0)

LSR# TF 7.1t

Tape T 750-1

Abstract: If this subroutine is entered at register Or, it will calculate cosine $\frac{\pi}{2}y$ by changing y to x so that cosine $\frac{\pi}{2}y = \text{sine } \frac{\pi}{2}x$ and evaluating sine $\frac{\pi}{2}x = ax - bx^3 + cx^5 - dx^7$, $-1 < x < 1$. If it is entered at register 7r, it will calculate sine $\frac{\pi}{2}x$. There will be an overflow at register 5r if $y=0$.

Upon entering the subroutine:

ac: x or y

Upon leaving the subroutine:

ac: sine $\frac{\pi}{2}x$ or cosine $\frac{\pi}{2}y$.

Temporary registers:

d - unused

1t - used to store x

Cosine →	00	ta 21r	Set return address	14	ad 26r	$-bx^2 - dx^4$
	01	ts 1t	y in 1t	15	mh 1t	$-bx + cx^3 - dx^5$
	02	op 4r	Is y negative? Yes	16	mh 1t	$-bx^2 + cx^4 - dx^6$
	03	cs 1t	No. -y in ac	17	sr* 1	$(-bx^2 + cx^4 - dx^6)2^{-1}$
2r →	04	ad 22r	$(1 - 2^{-15}) - y $ in ac	18	ad 27r	$(a - bx^2 + cx^4 - dx^6)2^{-1}$
	05	ad 23r	$1 - y = x$ in ac	19	mh 1t	$(ax - bx^3 + cx^5 - dx^7)2^{-1}$
	06	sp 8r	Possible overflow	20	sl 1	$ax - bx^3 + cx^5 - dx^7$
Sine →	07	ta 21r	Set return address (Or) (7r)	21	sp (0)	Return to main program
6r →	08	ts 1t	x in 1t	22	0.77777	$1 - 2^{-15}$
	09	mh 1t	x^2	23	0.00001	2^{-15}
	10	mh 24r	$-dx^2$	24	1.77560	-d
	11	ad 25r	c $-dx^2$	25	0.05055	c
	12	mh 1t	cx $-dx^3$	26	1.26521	-b
	13	mh 1t	cx $-dx^4$	27	0.62210	a x 2^{-1}